

## AT25XE011

### 1-Mbit, 1.65 V Minimum SPI Serial Flash Memory with Dual-Read Support

#### Features

- Single 1.65 V - 3.6 V Supply
- Serial Peripheral Interface (SPI) Compatible
  - Supports SPI Modes 0 and 3
  - Supports Dual Output Read
- 104 MHz Maximum Operating Frequency
  - Clock-to-Output (t<sub>v</sub>) of 6 ns
- Flexible, Optimized Erase Architecture for Code + Data Storage Applications
  - Small (256-Byte) Page Erase
  - Uniform 4-kByte Block Erase
  - Uniform 32-kByte Block Erase
  - Full Chip Erase
- Hardware Controlled Locking of Protected Sectors via WP Pin
- 128-byte, One-Time Programmable (OTP) Security Register
  - 64 bytes factory programmed with a unique identifier
  - 64 bytes user programmable
- Flexible Programming
  - Byte/Page Program (1 to 256 Bytes)
- Fast Program and Erase Times
  - 2 ms Typical Page Program (256 Bytes) Time
  - 50 ms Typical 4-kByte Block Erase Time
  - 400 ms Typical 32-kByte Block Erase Time
- Automatic Checking and Reporting of Erase/Program Failures
- Software Controlled Reset
- JEDEC Standard Manufacturer and Device ID Read Methodology
- Low Power Dissipation
  - 200 nA Ultra-Deep Power-Down current (Typical)
  - 4.5 µA Deep Power-Down Current (Typical)
  - 25 µA Standby current (Typical)
  - 3.5 mA Active Read Current (Typical)
- Endurance: 100,000 Program/Erase Cycles
- Data Retention: 20 Years
- Temperature Range: -10 °C to +85 °C (1.65 V to 3.6 V), -40 °C to +85 °C (1.7 V to 3.6 V)
- Industry Standard Green (Pb/Halide-free/RoHS Compliant) Package Options
  - 8-lead SOIC (150-mil)
  - 8-pad Ultra Thin DFN (2 x 3 x 0.6 mm)
  - 8-lead TSSOP Package

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## 1. Description

The AT25XE011 is a highly optimized ultra-low-energy serial interface Flash memory device designed for use in a wide variety of high-volume low-energy consumer and industrial applications. The AT25XE011 device has been optimized by design to meet the needs of today's connected applications in the Internet of Things market space.

The granular Page Erase and Block Erase architecture allows the memory space to be used much more efficiently supporting data storage and over the air updates. Key program code subroutines and data storage segments need to reside by themselves in their own separate erase regions, and with a granular architecture the wasted and unused memory space that occurs with large sectorized and large block erase flash memory devices can be greatly reduced.

The resulting improvement to software efficiency contributes to reduced CPU / MCU overheads that translate to further reduce the system energy usage levels. The device also contains a specialized OTP (One-Time Programmable) Security Register that can be used for purposes such as unique device serialization, system-level Electronic Serial Number (ESN) storage, locked key storage, etc. Specifically designed for use in many different systems, the AT25XE011 supports read, program, and erase operations with a wide supply voltage range of 1.65 V to 3.6 V. No separate voltage is required for programming and erasing.

## 2. Pin Descriptions and Pinouts

### 2.1 Descriptions

Table 1. Pin Descriptions

Symbol	Name and Function	Asserted State	Type
$\overline{\text{CS}}$	<b>CHIP SELECT:</b> Asserting the $\overline{\text{CS}}$ pin selects the device. When the $\overline{\text{CS}}$ pin is deasserted, the device is deselected and typically placed in standby mode (not Deep Power-Down mode), and the SO pin is in a high-impedance state. When the device is deselected, data is not accepted on the SI pin. A high-to-low transition on the $\overline{\text{CS}}$ pin is required to start an operation, and a low-to-high transition is required to end an operation. When ending an internally self-timed operation such as a program or erase cycle, the device does not enter the standby mode until the completion of the operation.	Low	Input
SCK	<b>SERIAL CLOCK:</b> This pin is used to provide a clock to the device and is used to control the flow of data to and from the device. Command, address, and input data present on the SI pin is always latched in on the rising edge of SCK, while output data on the SO pin is always clocked out on the falling edge of SCK.	-	Input
SI (I/O <sub>0</sub> )	<b>SERIAL INPUT:</b> The SI pin is used to shift data into the device. The SI pin is used for all data input including command and address sequences. Data on the SI pin is always latched in on the rising edge of SCK. With the Dual-Output Read commands, the SI Pin becomes an output pin (I/O <sub>0</sub> ) in conjunction with other pins to allow two bits of data on (I/O1-0) to be clocked out on every falling edge of SCK. To maintain consistency with the SPI nomenclature, the SI (I/O <sub>0</sub> ) pin is referenced as the SI pin unless specifically addressing the Dual-I/O modes in which case it is referenced as I/O0. Data present on the SI pin are ignored whenever the device is deselected ( $\overline{\text{CS}}$ is deasserted).	-	Input/Output
SO (I/O <sub>1</sub> )	<b>SERIAL OUTPUT:</b> The SO pin is used to shift data out from the device. Data on the SO pin is always clocked out on the falling edge of SCK. With the Dual-Output Read commands, the SO Pin remains an output pin (I/O <sub>1</sub> ) in conjunction with other pins to allow two bits of data on (I/O1-0) to be clocked out on every falling edge of SCK. To maintain consistency with the SPI nomenclature, the SO (I/O <sub>1</sub> ) pin is referenced as the SO pin unless specifically addressing the Dual-I/O modes in which case it is referenced as I/O1. The SO pin is in a high-impedance state whenever the device is deselected ( $\overline{\text{CS}}$ is deasserted).	-	Input/Output
$\overline{\text{WP}}$	<b>WRITE PROTECT:</b> The $\overline{\text{WP}}$ pin controls the hardware locking feature of the device. See <a href="#">"Protection Commands and Features" on page 20</a> for more details on protection features and the $\overline{\text{WP}}$ pin. The $\overline{\text{WP}}$ pin is internally pulled-high and can be left floating if hardware controlled protection is not used. However, it is recommended that the $\overline{\text{WP}}$ pin also be externally connected to VCC whenever possible.	Low	Input

Table 1. Pin Descriptions

$\overline{\text{HOLD}}$	<p><b>HOLD:</b></p> <p>The <math>\overline{\text{HOLD}}</math> pin is used to temporarily pause serial communication without deselecting or resetting the device. While the <math>\overline{\text{HOLD}}</math> pin is asserted, transitions on the SCK pin and data on the SI pin are ignored, and the SO pin is in a high-impedance state.</p> <p>The <math>\overline{\text{CS}}</math> pin must be asserted, and the SCK pin must be in the low state in order for a Hold condition to start. A Hold condition pauses serial communication only and does not have an effect on internally self-timed operations such as a program or erase cycle. See “Hold” on <a href="#">page 37</a> for additional details on the Hold operation.</p> <p>The <math>\overline{\text{HOLD}}</math> pin is internally pulled-high and can be left floating if the Hold function is not used. However, it is recommended that the <math>\overline{\text{HOLD}}</math> pin also be externally connected to <math>V_{\text{CC}}</math> whenever possible.</p>	Low	Input
$V_{\text{CC}}$	<p><b>DEVICE POWER SUPPLY:</b></p> <p>The <math>V_{\text{CC}}</math> pin is used to supply the source voltage to the device. Operations at invalid <math>V_{\text{CC}}</math> voltages can produce spurious results and must not be attempted.</p>	-	Power
GND	<p><b>GROUND:</b></p> <p>The ground reference for the power supply. GND must be connected to the system ground.</p>	-	Power

2.2 Pinouts

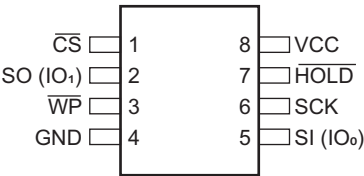


Figure 1. 8-SOIC Top View

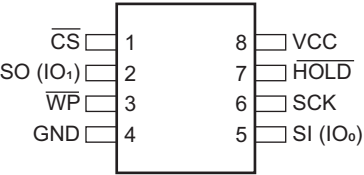


Figure 2. 8-TSSOP Top View

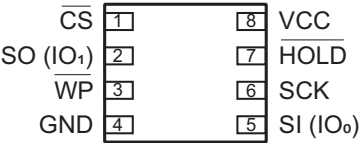


Figure 3. 8-UDFN (Top View)



### 3. Block Diagram

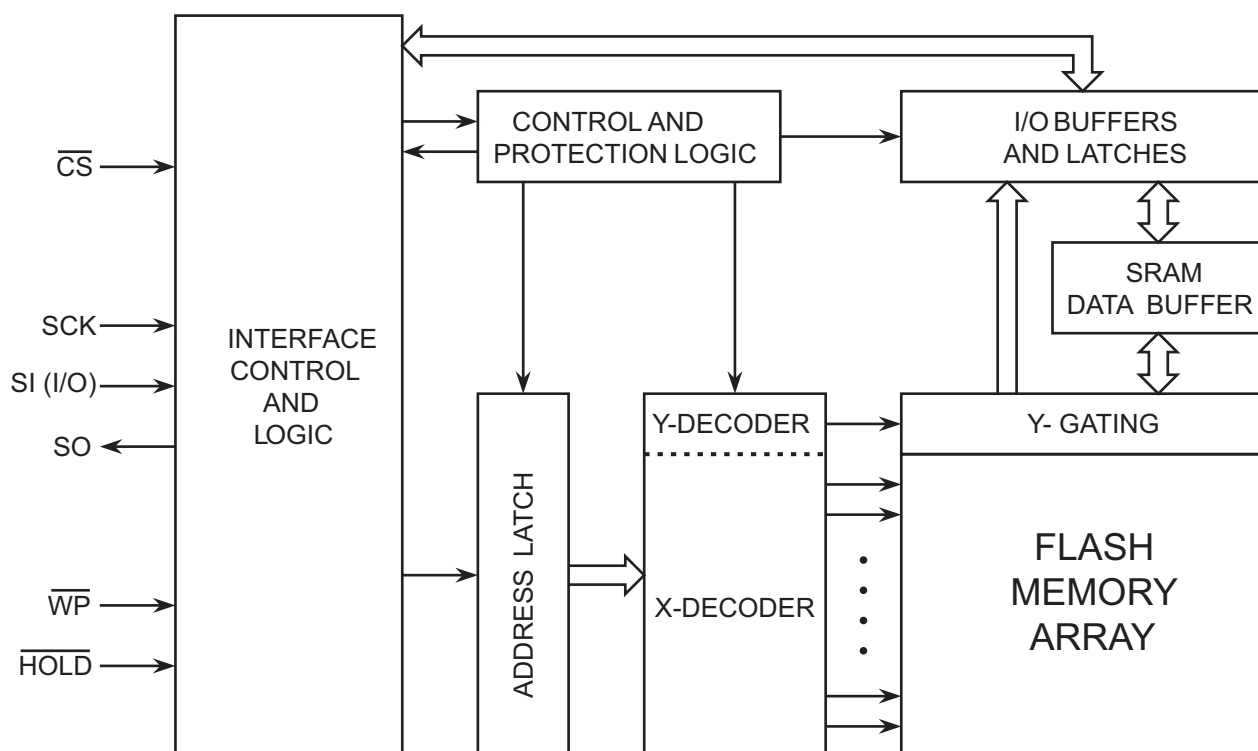


Figure 4. AT25XE011 Block Diagram

## 4. Memory Array

To provide the greatest flexibility, the memory array of the AT25XE011 can be erased in three levels of granularity, including a full-chip erase. The size of the erase blocks is optimized for both code and data storage applications, allowing both code and data segments to reside in their own erase regions. The Memory Architecture Diagram shows each erase level.

Block Erase Detail			Page Program Detail	
32 kB Block Erase (52h Command)	4 kB Block Erase (20h Command)	Block Address Range	1-256 Byte Page Program (02h Command)	Page Address Range
32 kB	4KB	01F000h - 01FFFFh	256 Bytes	01FF00h - 01FFFFh
	4KB	01E000h - 01EFFFh	256 Bytes	01FE00h - 01EFFFh
	4KB	01D000h - 01DFFFh	256 Bytes	01FD00h - 01DFFFh
	4KB	01C000h - 01CFFFh	256 Bytes	01FC00h - 01CFFFh
	4KB	01B000h - 01BFFFh	256 Bytes	01FB00h - 01BFFFh
	4KB	01A000h - 01AFFFh	256 Bytes	01FA00h - 01AFFFh
	4KB	019000h - 019FFFh	256 Bytes	01F900h - 01F9FFFh
	4KB	018000h - 018FFFh	256 Bytes	01F800h - 01F8FFFh
32 kB	4KB	017000h - 017FFFh	256 Bytes	01F700h - 01F7FFFh
	4KB	016000h - 016FFFh	256 Bytes	01F600h - 01F6FFFh
	4KB	015000h - 015FFFh	256 Bytes	01F500h - 01F5FFFh
	4KB	014000h - 014FFFh	256 Bytes	01F400h - 01F4FFFh
	4KB	013000h - 013FFFh	256 Bytes	01F300h - 01F3FFFh
	4KB	012000h - 012FFFh	256 Bytes	01F200h - 01F2FFFh
	4KB	011000h - 011FFFh	256 Bytes	01F100h - 01F1FFFh
	4KB	010000h - 010FFFh		
32 kB	4KB	00F000h - 00FFFFh		
	4KB	00E000h - 00EFFFh	256 Bytes	000E00h - 000EFFFh
	4KB	00D000h - 00DFFFh	256 Bytes	000D00h - 000DFFFh
	4KB	00C000h - 00CFFFh	256 Bytes	000C00h - 000CFFFh
	4KB	00B000h - 00BFFFh	256 Bytes	000B00h - 000BFFFh
	4KB	00A000h - 00AFFFh	256 Bytes	000A00h - 000AFFFh
	4KB	009000h - 009FFFh	256 Bytes	000900h - 0009FFFh
	4KB	008000h - 008FFFh	256 Bytes	000800h - 0008FFFh
32 kB	4KB	007000h - 007FFFh	256 Bytes	000700h - 0007FFFh
	4KB	006000h - 006FFFh	256 Bytes	000600h - 0006FFFh
	4KB	005000h - 005FFFh	256 Bytes	000500h - 0005FFFh
	4KB	004000h - 004FFFh	256 Bytes	000400h - 0004FFFh
	4KB	003000h - 003FFFh	256 Bytes	000300h - 0003FFFh
	4KB	002000h - 002FFFh	256 Bytes	000200h - 0002FFFh
	4KB	001000h - 001FFFh	256 Bytes	000100h - 0001FFFh
	4KB	000000h - 000FFFh	256 Bytes	000000h - 0000FFFh

Figure 5. Memory Architecture Diagram

## 5. Device Operation

The AT25XE011 is controlled by a set of instructions that are sent from a host controller, referred to as the SPI Master. The SPI Master communicates with the AT25XE011 through the SPI bus which is comprised of four signal lines: Chip Select ( $\overline{CS}$ ), Serial Clock (SCK), Serial Input (SI), and Serial Output (SO).

The SPI protocol defines a total of four modes of operation (mode 0, 1, 2, or 3) with each mode differing in respect to the SCK polarity and phase and how the polarity and phase control the flow of data on the SPI bus. The AT25XE011 supports the two most common modes, SPI Modes 0 and 3. The only difference between SPI Modes 0 and 3 is the polarity of the SCK signal when in the inactive state (when the SPI Master is in standby mode and not transferring any data). With SPI Modes 0 and 3, data is always latched in on the rising edge of SCK and always output on the falling edge of SCK.

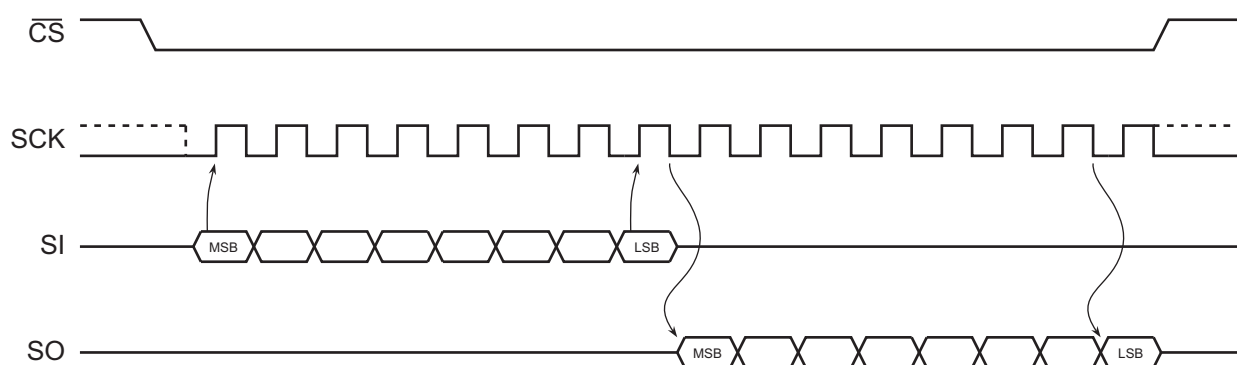


Figure 6. SPI Mode 0 and 3

### 5.1 Dual Output Read

The ATx features a Dual-Output Read mode that allow two bits of data to be clocked out of the device every clock cycle to improve throughput. To accomplish this, both the SI and SO pins are utilized as outputs for the transfer of data bytes. With the Dual-Output Read Array command, the SI pin becomes an output along with the SO pin.

## 6. Commands and Addressing

A valid instruction or operation must always be started by first asserting the  $\overline{CS}$  pin. After the pin has been asserted, the host controller must then clock out a valid 8-bit opcode on the SPI bus. Following the opcode, instruction dependent information such as address and data bytes would then be clocked out by the host controller. All opcode, address, and data bytes are transferred with the most-significant bit (MSB) first. An operation is ended by deasserting the  $\overline{CS}$  pin.

Opcodes not supported by the AT25XE011 are ignored by the device and no operation can start. The device continues to ignore any data presented on the SI pin until the start of the next operation ( $\overline{CS}$  pin being deasserted and then reasserted). Also, if the  $\overline{CS}$  pin is deasserted before complete opcode and address information is sent to the device, then no operation is performed and the device simply returns to the idle state and waits for the next operation.

Addressing of the device requires a total of three bytes of information to be sent, representing address bits A23-A0. Since the upper address limit of the AT25XE011 memory array is 01FFFFh, address bits A23-A17 are always ignored by the device.

**Table 2. Command Listing**

Command	Opcode	Clock Frequency	Address Bytes	Dummy Bytes	Data Bytes
<b>Read Commands</b>					
Read Array	0Bh    0000 1011	Up to 104 MHz	3	1	1+
	03h    0000 0011	Up to 33 MHz <sup>[1]</sup>	3	0	1+
Dual Output Read	3Bh    0011 1011	Up to 50 MHz	3	1	1+
<b>Program and Erase Commands</b>					
Page Erase	81h    1000 0001	Up to 104 MHz	3	0	0
Block Erase (4 kBytes)	20h    0010 0000	Up to 104 MHz	3	0	0
Block Erase (32 kBytes)	52h    0101 0010	Up to 104 MHz	3	0	0
	D8h    1101 1000	Up to 104 MHz	3	0	0
Chip Erase	60h    0110 0000	Up to 104 MHz	0	0	0
	C7h    1100 0111	Up to 104 MHz	0	0	0
Chip Erase (Legacy Command)	62h    0110 0010	Up to 104 MHz	0	0	0
Byte/Page Program (1 to 256 Bytes)	02h    0000 0010	Up to 104 MHz	3	0	1+
<b>Protection Commands</b>					
Write Enable	06h    0000 0110	Up to 104 MHz	0	0	0
Write Disable	04h    0000 0100	Up to 104 MHz	0	0	0
<b>Security Commands</b>					
Program OTP Security Register	9Bh    1001 1011	Up to 104 MHz	3	0	1+
Read OTP Security Register	77h    0111 0111	Up to 104 MHz	3	2	1+
<b>Status Register Commands</b>					
Read Status Register	05h    0000 0101	Up to 104 MHz	0	0	1+
Write Status Register Byte 1	01h    0000 0001	Up to 104 MHz	0	0	1
Write Status Register Byte 2	31h    0011 0001	Up to 104 MHz	0	0	1

Table 2. Command Listing

Miscellaneous Commands						
Reset	F0h	1111 0000	Up to 104 MHz	0	0	1(D0h)
Read Manufacturer and Device ID	9Fh	1001 1111	Up to 104 MHz	0	0	1 to 4
Read ID (Legacy Command)	15h	0001 0101	Up to 104 MHz	0	0	2
Deep Power-Down	B9h	1011 1001	Up to 104 MHz	0	0	0
Resume from Deep Power-Down	ABh	1010 1011	Up to 104 MHz	0	0	0
Ultra-Deep Power-Down	79h	0111 1001	Up to 104 MHz	0	0	0

1. Varies by voltage range. See [“AC Characteristics - Maximum Clock Frequencies” on page 40](#).

## 7. Read Commands

### 7.1 Read Array

The Read Array command can be used to sequentially read a continuous stream of data from the device by simply providing the clock signal once the initial starting address is specified. The device incorporates an internal address counter that automatically increments every clock cycle.

Two opcodes (0Bh and 03h) can be used for the Read Array command. The use of each opcode depends on the maximum clock frequency that is used to read data from the device. The 0Bh opcode can be used at any clock frequency up to the maximum specified by  $f_{CLK}$ , and the 03h opcode can be used for lower frequency read operations up to the maximum specified by  $f_{RDLF}$ .

To perform the Read Array operation, the  $\overline{CS}$  pin must first be asserted and the appropriate opcode (0Bh or 03h) must be clocked into the device. After the opcode has been clocked in, the three address bytes must be clocked in to specify the starting address location of the first byte to read within the memory array. Following the three address bytes, an additional dummy byte needs to be clocked into the device if the 0Bh opcode is used for the Read Array operation.

After the three address bytes (and the dummy byte if using opcode 0Bh) have been clocked in, additional clock cycles result in data being output on the SO pin. The data is always output with the MSB of a byte first. When the last byte (01FFFFh) of the memory array has been read, the device continues reading back at the beginning of the array (000000h). No delays are incurred when wrapping around from the end of the array to the beginning of the array.

Deasserting the  $\overline{CS}$  pin terminates the read operation and put the SO pin into high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require a full byte of data be read.

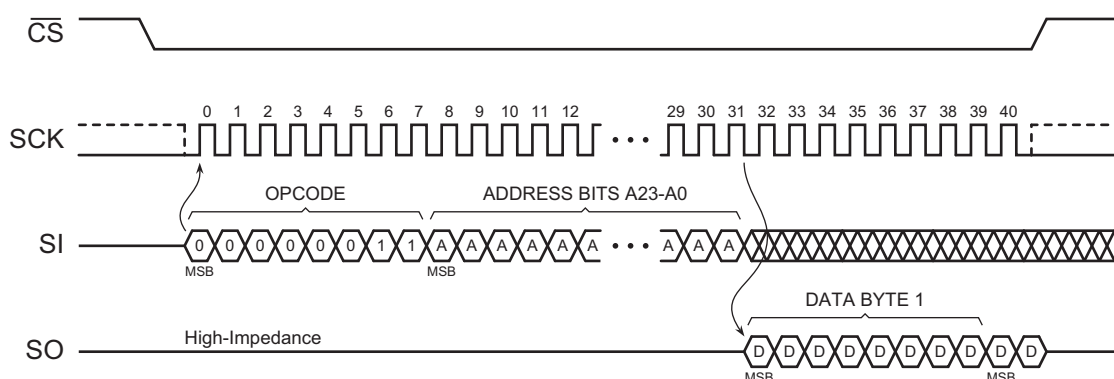


Figure 7. Read Array 03h Opcode

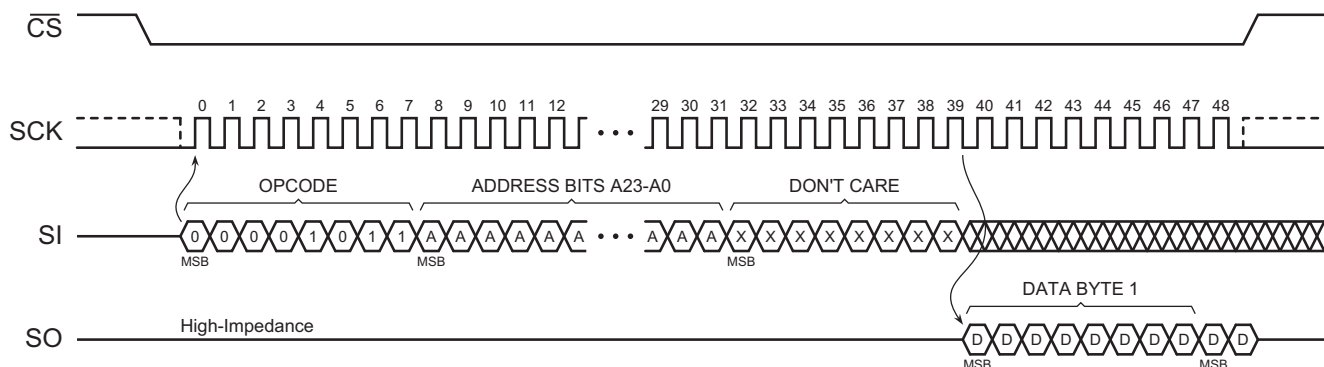


Figure 8. Read Array 0Bh Opcode

## 7.2 Dual-Output Read Array

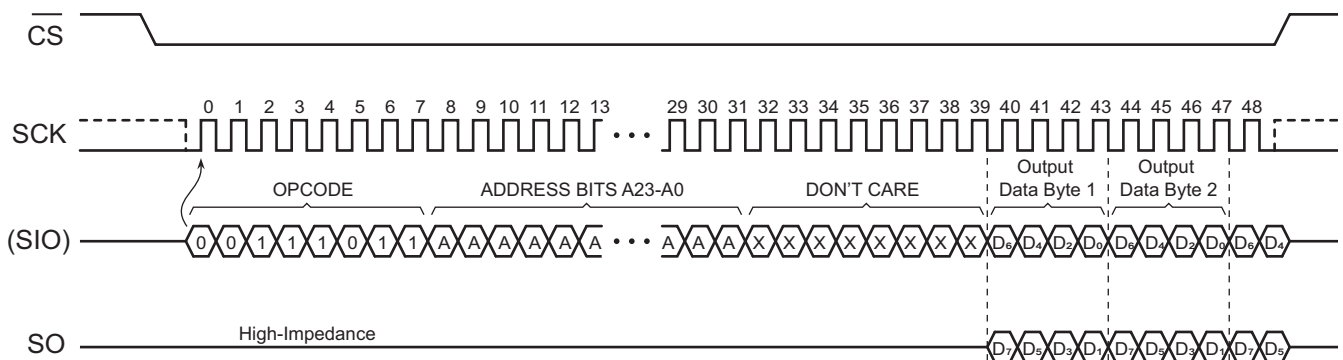
The Dual-Output Read Array command is similar to the standard Read Array command and can be used to sequentially read a continuous stream of data from the device by simply providing the clock signal once the initial starting address has been specified. Unlike the standard Read Array command, however, the Dual-Output Read Array command allows two bits of data to be clocked out of the device on every clock cycle, rather than just one.

The Dual-Output Read Array command can be used at any clock frequency, up to the maximum specified by  $f_{RDDO}$ . To perform the Dual-Output Read Array operation, the  $\overline{CS}$  pin must first be asserted and then the opcode 3Bh must be

clocked into the device. After the opcode has been clocked in, the three address bytes must be clocked in to specify the location of the first byte to read within the memory array. Following the three address bytes, a single dummy byte must also be clocked into the device.

After the three address bytes and the dummy byte have been clocked in, additional clock cycles result in data being output on both the SO and SIO pins. The data is always output with the MSB of a byte first and the MSB is always output on the SO pin. During the first clock cycle, bit seven of the first data byte is output on the SO pin, while bit six of the same data byte is output on the SIO pin. During the next clock cycle, bits five and four of the first data byte are output on the SO

and SIO pins, respectively. The sequence continues with each byte of data being output after every four clock cycles. When the last byte (01FFFFh) of the memory array has been read, the device continues reading from the beginning of the array (000000h). No delays are incurred when wrapping around from the end of the array to the beginning of the array. Deasserting the  $\overline{\text{CS}}$  pin terminates the read operation and put the SO and SIO pins into a high-impedance state. The  $\overline{\text{CS}}$  pin can be deasserted at any time and does not require that a full byte of data be read.



### Figure 9. Dual Output Read Array

## 8. Program and Erase Commands

### 8.1 Byte/Page Program

The Byte/Page Program command allows anywhere from a single byte of data to 256 bytes of data to be programmed into previously erased memory locations. An erased memory location is one that has all eight bits set to the logical 1 state (a byte value of FFh). Before a Byte/Page Program command can be started, the Write Enable command must have been previously issued to the device (see [“Write Enable” on page 20](#)) to set the Write Enable Latch (WEL) bit of the Status Register to a logical 1 state.

To perform a Byte/Page Program command, an opcode of 02h must be clocked into the device followed by the three address bytes denoting the first byte location of the memory array to begin programming at. After the address bytes have been clocked in, data can then be clocked into the device and are stored in an internal buffer.

If the starting memory address denoted by A23-A0 does not fall on an even 256-byte page boundary (A7-A0 are not all 0), then special circumstances regarding which memory locations to be programmed apply. In this situation, any data that is sent to the device that goes beyond the end of the page wraps around back to the beginning of the same page. For example, if the starting address denoted by A23-A0 is 0000FEh, and three bytes of data are sent to the device, then the first two bytes of data are programmed at addresses 0000FEh and 0000FFh while the last byte of data is programmed at address 000000h. The remaining bytes in the page (addresses 000001h through 0000FDh) are not programmed and remains in the erased state (FFh). Also, if more than 256 bytes of data are sent to the device, then only the last 256 bytes sent are latched into the internal buffer.

When the  $\overline{\text{CS}}$  pin is deasserted, the device takes the data stored in the internal buffer and program it into the appropriate memory array locations based on the starting address specified by A23-A0 and the number of data bytes sent to the device. If less than 256 bytes of data were sent to the device, then the remaining bytes within the page are not programmed and remain in the erased state (FFh). The programming of the data bytes is internally self-timed and must take place in a time of  $t_{\text{BP}}$  or  $t_{\text{PP}}$  if only programming a single byte.

The three address bytes and at least one complete byte of data must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device aborts the operation and no data are programmed into the memory array. Also, if the memory is in the protected state (see [“Block Protection” on page 21](#)), then the Byte/Page Program command is not executed, and the device returns to the idle state once the  $\overline{\text{CS}}$  pin has been deasserted. The WEL bit in the Status Register is reset back to the logical 0 state if the program cycle aborts due to an incomplete address being sent, an incomplete byte of data being sent, the  $\overline{\text{CS}}$  pin being deasserted on uneven byte boundaries, or because the memory location to be programmed is protected.

While the device is programming, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{\text{BP}}$  or  $t_{\text{PP}}$  time to determine if the data bytes have finished programming. At some point before the program cycle completes, the WEL bit in the Status Register is reset back to the logical 0 state.

The device also incorporates an intelligent programming algorithm that can detect when a byte location fails to program properly. If a programming error arises, it is indicated by the EPE bit in the Status Register.



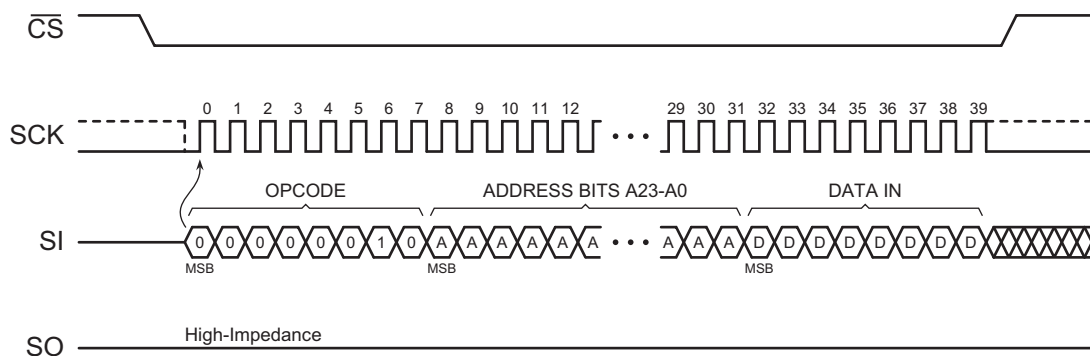


Figure 10. Byte Program

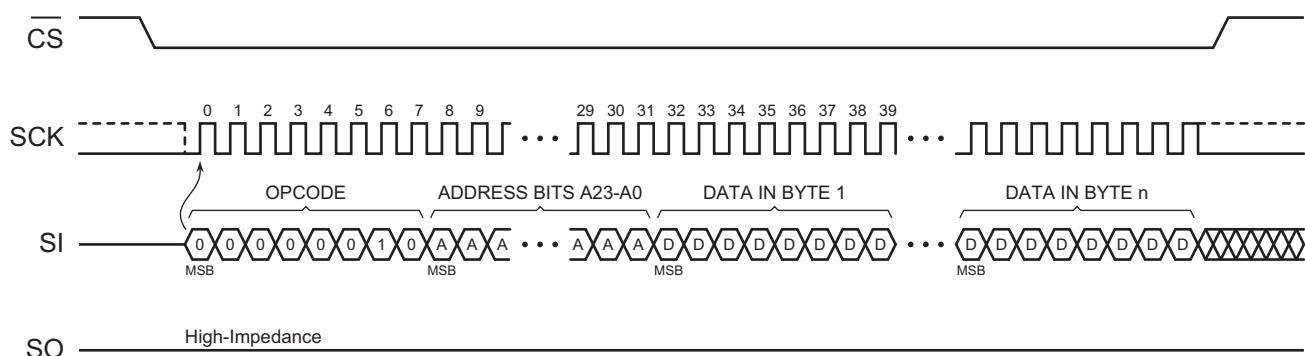


Figure 11. Page Program

## 8.2 Page Erase

Page Erase for 1 Mbit, 512 Pages [nine (9) page address bits, PA<8:0>] of 256 Bytes each.

The Page Erase command can be used to individually erase any page in the main memory array. The Main Memory Byte/Page Program command can be utilized at a later time.

To perform a Page Erase with the standard page size (256 bytes), an opcode of 81h must be clocked into the device followed by three address bytes comprised of:

Byte 0: 81h the page erase command code

Byte 1: XXXX XXX, PA8; which is seven (7) dummy bits, and one (1) page address bit PA8  
Byte 2: PA<7:0>; which is eight (8) page address bits

Byte 3: XXXX XXXX; which is eight (8) dummy bits

When a low-to-high transition occurs on the  $\overline{CS}$  pin, the device erases the selected page (the erased state is a Logic 1). The erase operation is internally self-timed and must take place in a maximum time of  $t_{PE}$ . During this time, the RDY/BUSY bit in the Status Register indicates that the device is busy.

The device also incorporates an intelligent erase algorithm that can detect when a byte location fails to erase properly. If an erase error arises, it is indicated by the EPE bit in the Status Register.

### 8.3 Block Erase

A block of 4 or 32 kBytes can be erased (all bits set to the logical 1 state) in a single operation by using one of three different opcodes for the Block Erase command. An opcode of 20h is used for a 4-kByte erase, and an opcode of 52h or D8h is used for a 32-kByte erase. Before a Block Erase command can be started, the Write Enable command must have been previously issued to the device to set the WEL bit of the Status Register to a logical 1 state.

To perform a Block Erase, the  $\overline{\text{CS}}$  pin must first be asserted and the appropriate opcode (20h, 52h, or D8h) must be clocked into the device. After the opcode has been clocked in, the three address bytes specifying an address within the 4- or 32-kByte block to be erased must be clocked in. Any additional data clocked into the device is ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the device erases the appropriate block. The erasing of the block is internally self-timed and must take place in a time of  $t_{\text{BLKE}}$ .

Since the Block Erase command erases a region of bytes, the lower order address bits do not need to be decoded by the device. Therefore, for a 4-kByte erase, address bits A11-A0 are ignored by the device and their values can be either a logical 1 or "0". For a 32-kByte erase, address bits A14-A0 are ignored by the device. Despite the lower order address bits not being decoded by the device, the complete three address bytes must still be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, the device aborts the operation and no erase operation are performed.

If the memory is in the protected state, then the Block Erase command is not be executed, and the device returns to the idle state once the  $\overline{\text{CS}}$  pin has been deasserted.

The WEL bit in the Status Register is reset back to the logical 0 state if the erase cycle aborts due to an incomplete address being sent, the  $\overline{\text{CS}}$  pin being deasserted on uneven byte boundaries, or because a memory location within the region to be erased is protected.

While the device is executing a successful erase cycle, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{\text{BLKE}}$  time to determine if the device has finished erasing. At some point before the erase cycle completes, the WEL bit in the Status Register is reset back to the logical 0 state.

The device also incorporates an intelligent erase algorithm that can detect when a byte location fails to erase properly. If an erase error occurs, it is indicated by the EPE bit in the Status Register.

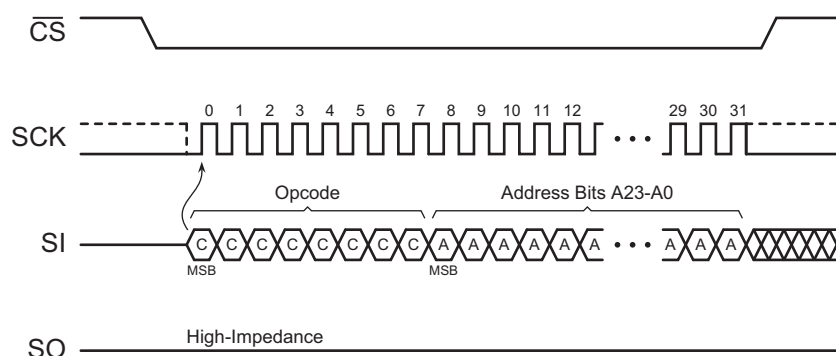


Figure 12. Block Erase

## 8.4 Chip Erase

The entire memory array can be erased in a single operation by using the Chip Erase command. Before a Chip Erase command can be started, the Write Enable command must have been previously issued to the device to set the WEL bit of the Status Register to a logical 1 state.

Three opcodes (60h, 62h, and C7h) can be used for the Chip Erase command. There is no difference in device functionality when utilizing the three opcodes, so they can be used interchangeably. To perform a Chip Erase, one of the three opcodes must be clocked into the device. Since the entire memory array is to be erased, no address bytes need to be clocked into the device, and any data clocked in after the opcode is ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the device erases the entire memory array. The erasing of the device is internally self-timed and must take place in a time of  $t_{\text{CHPE}}$ .

The complete opcode must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, no erase is performed. Also, if the memory array is in the protected state, then the Chip Erase command is not executed, and the device returns to the idle state once the  $\overline{\text{CS}}$  pin has been deasserted. The WEL bit in the Status Register is reset back to the logical 0 state if the  $\overline{\text{CS}}$  pin is deasserted on uneven byte boundaries or if the memory is in the protected state.

While the device is executing a successful erase cycle, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{\text{CHPE}}$  time to determine if the device has finished erasing. At some point before the erase cycle completes, the WEL bit in the Status Register is reset back to the logical 0 state.

The device also incorporates an intelligent erase algorithm that can detect when a byte location fails to erase properly. If an erase error occurs, it is indicated by the EPE bit in the Status Register.

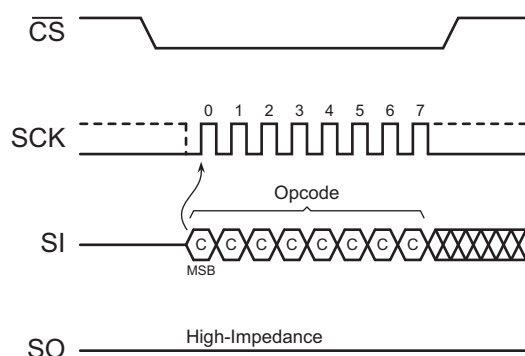


Figure 13. Chip Erase

## 9. Protection Commands and Features

### 9.1 Write Enable

The Write Enable command is used to set the Write Enable Latch (WEL) bit in the Status Register to a logical 1 state. The WEL bit must be set before a Byte/Page Program, erase, Program OTP Security Register, or Write Status Register command can be executed. This makes the issuance of these commands a two step process, thereby reducing the chances of a command being accidentally or erroneously executed. If the WEL bit in the Status Register is not set prior to the issuance of one of these commands, then the command is not executed.

To issue the Write Enable command, the  $\overline{\text{CS}}$  pin must first be asserted and the opcode of 06h must be clocked into the device. No address bytes need to be clocked into the device, and any data clocked in after the opcode is ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the WEL bit in the Status Register is set to a logical 1. The complete opcode must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, the device aborts the operation and the state of the WEL bit does not change.

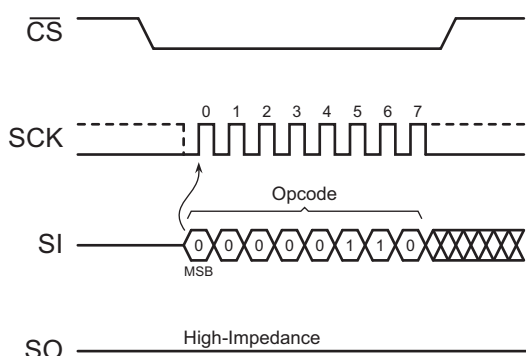


Figure 14. Write Enable

### 9.2 Write Disable

The Write Disable command is used to reset the Write Enable Latch (WEL) bit in the Status Register to the logical 0 state. With the WEL bit reset, all Byte/Page Program, erase, Program OTP Security Register, and Write Status Register commands are not executed. Other conditions can also cause the WEL bit to be reset; see the WEL bit section of the Status Register description for more details.

To issue the Write Disable command, the  $\overline{\text{CS}}$  pin must first be asserted and the opcode of 04h must be clocked into the device. No address bytes need to be clocked into the device, and any data clocked in after the opcode is ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the WEL bit in the Status Register is reset to a logical 0. The complete opcode must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, the device aborts the operation and the state of the WEL bit does not change.

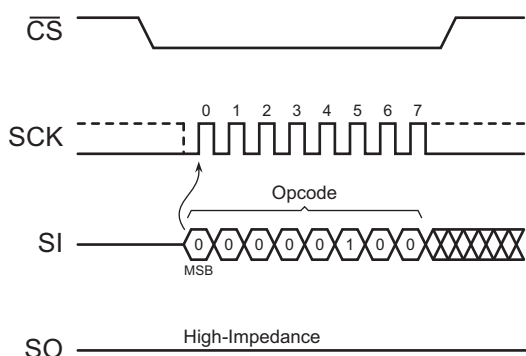


Figure 15. Write Disable

### 9.3 Block Protection

The device can be software protected against erroneous or malicious program or erase operations by utilizing the Block Protection feature of the device. Block Protection can be enabled or disabled by using the Write Status Register command to change the value of the Block Protection (BP0) bit in the Status Register. The following table outlines the two states of the BP0 bit and the associated protection area.

**Table 3. Memory Array Protection**

Protection Level	BP0	Protected Memory Address
None	0	None
Full Memory	1	00000h - 01FFFFh

When the BP0 bit of the Status Register is in the logical 1 state, the entire memory array is protected against program or erase operations. Any attempts to send a Byte/Page Program command, a Block Erase command, or a Chip Erase command are ignored by the device.

As a safeguard against accidental or erroneous protecting or unprotecting of the memory array, the BP0 bit itself can be locked from updates by using the  $\overline{WP}$  pin and the BPL (Block Protection Locked) bit of the Status Register (see [“Protected States and the Write Protect Pin” on page 22](#) for more details).

The BP0 bit of the Status Register is a nonvolatile bit; therefore, the BP0 bit retains its state even after the device has been power cycled. Care must be taken to ensure that BP0 is in the logical 1 state before powering down for those applications that wish to have the memory array fully protected upon power up. The default state for BP0 when shipped from Renesas Electronics is 0.

## 9.4 Protected States and the Write Protect Pin

The  $\overline{\text{WP}}$  pin is not linked to the memory array itself and has no direct effect on the protection status of the memory array. Instead, the  $\overline{\text{WP}}$  pin, in conjunction with the BPL (Block Protection Locked) bit in the Status Register, is used to control the hardware locking mechanism of the device. For hardware locking to be active, two conditions must be met—the  $\overline{\text{WP}}$  pin must be asserted and the BPL bit must be in the logical 1 state.

When hardware locking is active, the Block Protection (BP0) bit is locked and the BPL bit itself is also locked. Therefore, if the memory array is protected, it is locked in the protected state, and if the memory array is unprotected, it is locked in the unprotected state. These states cannot be changed as long as hardware locking is active, so the Write Status Register command are ignored. In order to modify the protection status of the memory array, the  $\overline{\text{WP}}$  pin must first be deasserted, and the BPL bit in the Status Register must be reset back to the logical 0 state using the Write Status Register command.

If the  $\overline{\text{WP}}$  pin is permanently connected to GND, then once the BPL bit is set to a logical 1, the only way to reset the bit back to the logical 0 state is to power-cycle the device. This allows a system to power-up with all sectors software protected but not hardware locked. Therefore, sectors can be unprotected and protected as needed and then hardware locked at a later time by simply setting the BPL bit in the Status Register.

When the  $\overline{\text{WP}}$  pin is deasserted, or if the  $\overline{\text{WP}}$  pin is permanently connected to V<sub>CC</sub>, the BPL bit in the Status Register can be set to a logical 1, but doing so does not lock the BP0 bit.

Table 4 details the various protection and locking states of the device.

**Table 4. Hardware and Software Locking**

$\overline{\text{WP}}$	BPL	Locking	BPL Change Allowed	BP0 and Protection Status
0	0		Can be modified from 0 to 1	BP0 bit unlocked and modifiable using the Write Status Register command. Memory array can be protected and unprotected freely.
0	1	Hardware Locked	Locked	BP0 bit locked in current state. The Write Status Register command has no affect. Memory array is locked in current protected or unprotected state.
1	0		Can be modified from 0 to 1	BP0 bit unlocked and modifiable using the Write Status Register command. Memory array can be protected and unprotected freely.
1	1		Can be modified from 1 to 0	BP0 bit unlocked and modifiable using the Write Status Register command. Memory array can be protected and unprotected freely.

## 10. Security Commands

### 10.1 Program OTP Security Register

The device contains a specialized OTP (One-Time Programmable) Security Register that can be used for purposes such as unique device serialization, system-level Electronic Serial Number (ESN) storage, locked key storage, etc. The OTP Security Register is independent of the main Flash memory array and is comprised of a total of 128 bytes of memory divided into two portions. The first 64 bytes (byte locations 0 through 63) of the OTP Security Register are allocated as a one-time user-programmable space. Once these 64 bytes have been programmed, they cannot be erased or reprogrammed. The remaining 64 bytes of the OTP Security Register (byte locations 64 through 127) are factory programmed and contains a unique value for each device. The factory programmed data is fixed and cannot be changed.

**Table 5. OTP Security Register**

Security Register Byte Number									
0	1		62	63	64	65		126	127
One-Time User Programmable					Factory Programmed by Renesas Electronics				

The user-programmable portion of the OTP Security Register does not need to be erased before it is programmed. Also, the Program OTP Security Register command operates on the entire 64-byte user-programmable portion of the OTP Security Register at one time. Once the user-programmable space has been programmed with any number of bytes, the user-programmable space cannot be programmed again; therefore, it is not possible to only program the first two bytes of the register and then program the remaining 62 bytes at a later time.

Before the Program OTP Security Register command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a logical 1. To program the OTP Security Register, the  $\overline{\text{CS}}$  pin must first be asserted and an opcode of 9Bh must be clocked into the device followed by the three address bytes denoting the first byte location of the OTP Security Register to begin programming at. Since the size of the user-programmable portion of the OTP Security Register is 64 bytes, the upper order address bits do not need to be decoded by the device. Therefore, address bits A23-A6 are ignored by the device and their values can be either a logical 1 or 0. After the address bytes have been clocked in, data can then be clocked into the device and is stored in the internal buffer.

If the starting memory address denoted by A23-A0 does not start at the beginning of the OTP Security Register memory space (A5-A0 are not all 0), then special circumstances regarding which OTP Security Register locations to be programmed apply. In this situation, any data that is sent to the device that goes beyond the end of the 64-byte user-programmable space wraps around back to the beginning of the OTP Security Register. For example, if the starting address denoted by A23-A0 is 00003Eh, and three bytes of data are sent to the device, then the first two bytes of data are programmed at OTP Security Register addresses 00003Eh and 00003Fh while the last byte of data is programmed at address 000000h. The remaining bytes in the OTP Security Register (addresses 000001h through 00003Dh) are not programmed and remain in the erased state (FFh). Also, if more than 64 bytes of data are sent to the device, then only the last 64 bytes sent are latched into the internal buffer.

When the  $\overline{\text{CS}}$  pin is deasserted, the device takes the data stored in the internal buffer and program it into the appropriate OTP Security Register locations based on the starting address specified by A23-A0 and the number of data bytes sent to the device. If less than 64 bytes of data were sent to the device, then the remaining bytes within the OTP Security Register are not programmed and remain in the erased state (FFh). The programming of the data bytes is internally self-timed and must take place in a time of  $t_{\text{OTPP}}$ .

The three address bytes and at least one complete byte of data must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device aborts the operation and the user-programmable portion of the OTP Security Register is not programmed. The WEL bit in the Status Register is reset back to the logical 0 state if the OTP Security Register program cycle

aborts due to an incomplete address being sent, an incomplete byte of data being sent, the  $\overline{CS}$  pin being deasserted on uneven byte boundaries, or because the user-programmable portion of the OTP Security Register was previously programmed.

While the device is programming the OTP Security Register, the Status Register can be read and indicates that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the  $t_{OTP}$  time to determine if the data bytes have finished programming. At some point before the OTP Security Register programming completes, the WEL bit in the Status Register is reset back to the logical 0 state.

If the device is powered-down during the OTP Security Register program cycle, then the contents of the 64-byte user programmable portion of the OTP Security Register cannot be guaranteed and cannot be programmed again.

The Program OTP Security Register command utilizes the internal 256-buffer for processing. Therefore, the contents of the buffer is altered from its previous state when this command is issued.

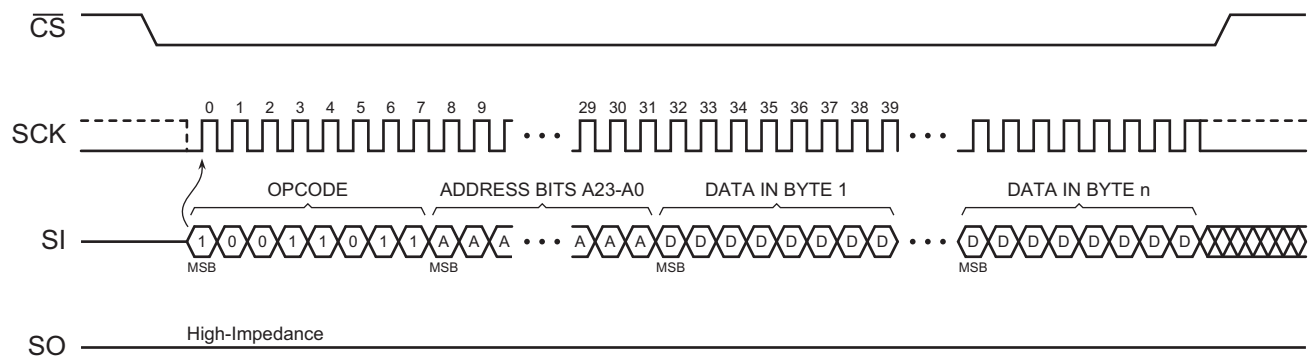


Figure 16. Program OTP Security Register

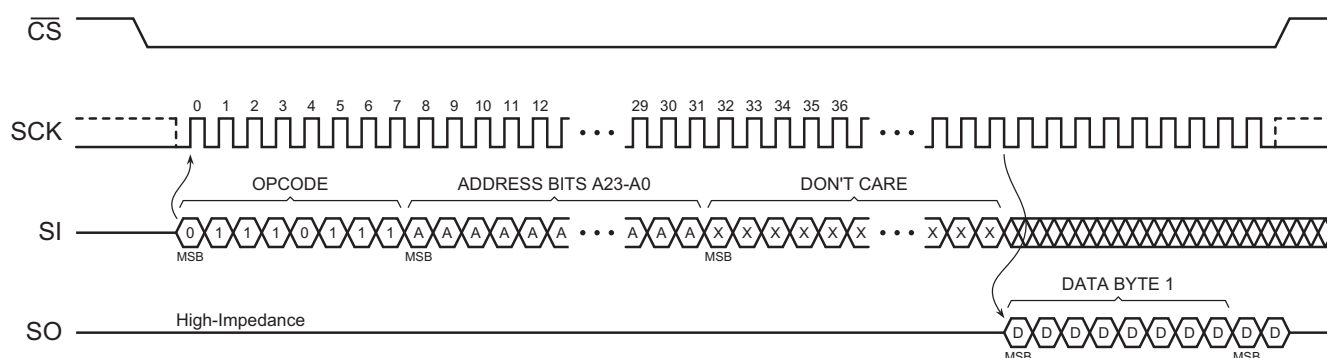


## 10.2 Read OTP Security Register

The OTP Security Register can be sequentially read in a similar fashion to the Read Array operation up to the maximum clock frequency specified by  $f_{CLK}$ . To read the OTP Security Register, the  $\overline{CS}$  pin must first be asserted and the opcode of 77h must be clocked into the device. After the opcode has been clocked in, the three address bytes must be clocked in to specify the starting address location of the first byte to read within the OTP Security Register. Following the three address bytes, two dummy bytes must be clocked into the device before data can be output.

After the three address bytes and the dummy bytes have been clocked in, additional clock cycles result in OTP Security Register data being output on the SO pin. When the last byte (00007Fh) of the OTP Security Register has been read, the device continues reading back at the beginning of the register (000000h). No delays are incurred when wrapping around from the end of the register to the beginning of the register.

Deasserting the  $\overline{\text{CS}}$  pin terminates the read operation and put the SO pin into a high-impedance state. The  $\overline{\text{CS}}$  pin can be deasserted at any time and does not require that a full byte of data be read.



### Figure 17. Read OTP Security Register

## 11. Status Register Commands

### 11.1 Read Status Register

The Status Register can be read to determine the device's ready/busy status, as well as the status of many other functions such as Hardware Locking and Block Protection. The Status Register can be read at any time, including during an internally self-timed program or erase operation. The Status Register consists of two bytes.

To read the Status Register, the  $\overline{CS}$  pin must first be asserted and the opcode of 05h must be clocked into the device. After the opcode has been clocked in, the device begins outputting Status Register data on the SO pin during every subsequent clock cycle. After the last bit (bit 0) of Status Register Byte 1 has been clocked out, the first bit (bit 7) of Status Register Byte 2 is clocked out. After the last bit (bit 0) of Status Register Byte 2 has been clocked out, the sequence repeats itself, starting again with bit 7 of Status Register Byte 1, as long as the  $\overline{CS}$  pin remains asserted and the clock pin is being pulsed. The data in the Status Register is constantly being updated, so each repeating sequence outputs new data.

Deasserting the  $\overline{CS}$  pin terminates the Read Status Register operation and put the SO pin into a high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require that a full byte of data be read.

**Table 6. Status Register Format**

Bit <sup>[1]</sup>	Name		Type <sup>[2]</sup>	Description	
7	BPL	Block Protection Locked	R/W	0	BP0 bit unlocked (default).
				1	BP0 bit locked in current state when $\overline{WP}$ asserted.
6	RES	Reserved for future use	R	0	Reserved for future use.
5	EPE	Erase/Program Error	R	0	Erase or program operation was successful.
				1	Erase or program error detected.
4	WPP	Write Protect ( $\overline{WP}$ ) Pin Status	R	0	$\overline{WP}$ is asserted.
				1	$\overline{WP}$ is deasserted.
3	RES	Reserved for future use	R	0	Reserved for future use.
2	BP0	Block Protection	R/W	0	Entire memory array is unprotected.
				1	Entire memory array is protected.
1	WEL	Write Enable Latch Status	R	0	Device is not write enabled (default).
				1	Device is write enabled.
0	RDY/BSY	Ready/Busy Status	R	0	Device is ready.
				1	Device is busy with an internal operation.

1. Only bits 7 and 2 of the Status Register can be modified when using the Write Status Register command.

2. R/W = Readable and Writeable; R = Readable only.

### 11.1.1 BPL Bit

The BPL bit is used to control whether the Block Protection (BP0) bit can be modified or not. When the BPL bit is in the logical 1 state and the  $\overline{WP}$  pin is asserted, the BP0 bit is locked and cannot be modified. The memory array is locked in the current protected or unprotected state.

When the BPL bit is in the logical 0 state, the BP0 bit is unlocked and can be modified. The BPL bit defaults to the logical 0 state after device power-up.

The BPL bit can be modified freely whenever the  $\overline{WP}$  pin is deasserted. However, if the  $\overline{WP}$  pin is asserted, then the BPL bit can only be changed from a logical 0 (BP0 bit unlocked) to a logical 1 (BP0 bit locked). In order to reset the BPL bit back to a logical 0 using the Write Status Register command, the  $\overline{WP}$  pin has to first be deasserted.

The BPL and BP0 bits are the only bits of the Status Register that can be user-modified through the Write Status Register command.

### 11.1.2 EPE Bit

The EPE bit indicates whether the last erase or program operation completed successfully or not. If at least one byte during the erase or program operation did not erase or program properly, then the EPE bit is set to the logical 1 state. The EPE bit is not set if an erase or program operation aborts for any reason such as an attempt to erase or program the memory when it is protected or if the WEL bit is not set prior to an erase or program operation. The EPE bit is updated after every erase and program operation.

### 11.1.3 WPP Bit

The WPP bit can be read to determine if the WP pin has been asserted or not.

### 11.1.4 BP0 Bit

The BP0 bits provides feedback on the software protection status for the device. Also, the BP0 bit can also be modified to change the state of the software protection to allow the entire memory array to be protected or unprotected. When the BP0 bit is in the logical 0 state, then the entire memory array is unprotected. When the BP0 bit is in the logical 1 state, then the entire memory array is protected against program and erase operations.

### 11.1.5 WEL Bit

The WEL bit indicates the current status of the internal Write Enable Latch. When the WEL bit is in the logical 0 state, the device does not accept any Byte/Page Program, erase, Program OTP Security Register, or Write Status Register commands. The WEL bit defaults to the logical 0 state after a device power-up or reset operation. Also, the WEL bit is reset to the logical 0 state automatically under the following conditions:

- Write Disable operation completes successfully
- Write Status Register operation completes successfully or aborts
- Program OTP Security Register operation completes successfully or aborts
- Byte/Page Program operation completes successfully or aborts
- Block Erase operation completes successfully or aborts
- Chip Erase operation completes successfully or aborts
- Hold condition aborts

If the WEL bit is in the logical 1 state, it is not reset to a logical 0 if an operation aborts due to an incomplete or unrecognized opcode being clocked into the device before the  $\overline{CS}$  pin is deasserted. In order for the WEL bit to be reset when an operation aborts prematurely, the entire opcode for a Byte/Page Program, erase, Program OTP Security Register, or Write Status Register command must have been clocked into the device

### 11.1.6 RDY/BSY Bit

The RDY/BSY bit is used to determine whether or not an internal operation, such as a program or erase, is in progress. To poll the RDY/BSY bit to detect the completion of a program or erase cycle, new Status Register data must be continually clocked out of the device until the state of the RDY/BSY bit changes from a logical 1 to a logical 0. Note that the RDY/BSY bit can be read either from Status Register Byte 1 or from Status Register Byte 2.

### 11.1.7 RSTE Bit

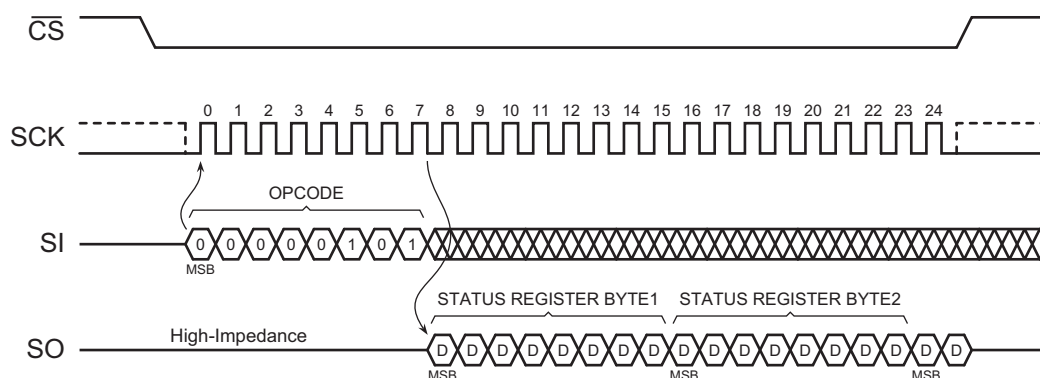
The RSTE bit is used to enable or disable the Reset command. When the RSTE bit is in the Logical 0 state (the default state after power-up), the Reset command is disabled and any attempts to reset the device using the Reset command are ignored. When the RSTE bit is in the Logical 1 state, the Reset command is enabled.

The RSTE bit retains its state as long as power is applied to the device. Once set to the Logical 1 state, the RSTE bit remains in that state until it is modified using the Write Status Register Byte 2 command or until the device has been power cycled. The Reset command itself does not change the state of the RSTE bit.

**Table 7. Status Register Format - Byte 2**

Bit <sup>[1]</sup>	Name		Type <sup>[2]</sup>	Description	
7	RES	Reserved for future use	R	0	Reserved for future use
6	RES	Reserved for future use	R	0	Reserved for future use
5	RES	Reserved for future use	R	0	Reserved for future use
4	RSTE	Reset Enabled	R/W	0	Reset command is disabled (default)
				1	Reset command is enabled
3	RES	Reserved for future use	R	0	Reserved for future use
2	RES	Reserved for future use	R	0	Reserved for future use
1	RES	Reserved for future use	R	0	Reserved for future use
0	RDY/BSY	Ready/Busy Status	R	0	Device is ready
				1	Device is busy with an internal operation

1. Only bits 4 and 3 of Status Register Byte 2 are modified when using the Write Status Register Byte 2 command.
2. R/W = Readable and Writeable; R = Readable only.



**Figure 18. Read Status Register**

## 11.2 Write Status Register

The Write Status Register command is used to modify the BPL bit and the BP0 bit of the Status Register. Before the Write Status Register command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a logical 1.

To issue the Write Status Register command, the  $\overline{\text{CS}}$  pin must first be asserted and the opcode of 01h must be clocked into the device followed by one byte of data. The one byte of data consists of the BPL bit value, four don't care bits, the BP0 bit value, and two additional don't care bits (see Table 8). Any additional data bytes that are sent to the device are ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the BPL bit and the BP0 bit in the Status Register is modified, and the WEL bit in the Status Register is reset back to a logical 0. The value of BP0 and the state of the BPL bit and the  $\overline{\text{WP}}$  pin before the Write Status Register command was executed (the prior state of the BPL bit and the state of the  $\overline{\text{WP}}$  pin when the  $\overline{\text{CS}}$  pin is deasserted) determines whether or not software protection is changed. See “Protected States and the Write Protect Pin” on page 22 for more details.

The complete one byte of data must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device aborts the operation, the state of the BPL and BP0 bits do not change, memory protection status does not change, and the WEL bit in the Status Register is reset back to the logical 0 state.

If the  $\overline{\text{WP}}$  pin is asserted, then the BPL bit can only be set to a logical 1. If an attempt is made to reset the BPL bit to a logical 0 while the  $\overline{\text{WP}}$  pin is asserted, then the Write Status Register Byte command is ignored, and the WEL bit in the Status Register is reset back to the logical 0 state. In order to reset the BPL bit to a logical 0, the  $\overline{\text{WP}}$  pin must be deasserted.

Table 8. Write Status Register Format

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
BPL	X	X	X	X	BP0	X	X

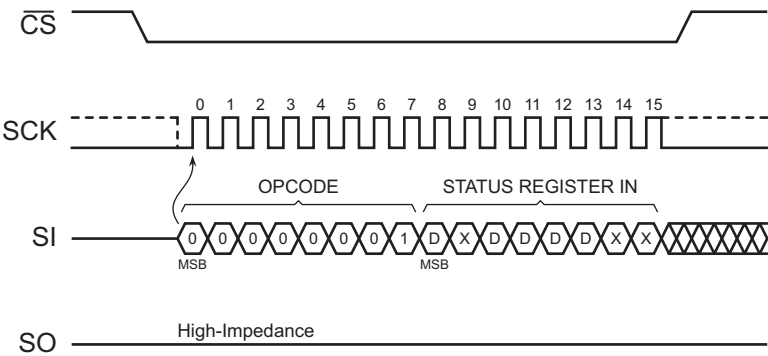


Figure 19. Write Status Register

### 11.3 Write Status Register Byte 2

The Write Status Register Byte 2 command is used to modify the RSTE. Using the Write Status Register Byte 2 command is the only way to modify the RSTE in the Status Register during normal device operation. Before the Write Status Register Byte 2 command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a Logical 1.

To issue the Write Status Register Byte 2 command, the  $\overline{\text{CS}}$  pin must first be asserted and then the opcode 31h must be clocked into the device followed by one byte of data. The one byte of data consists of three don't-care bits, the RSTE bit value, and four additional don't-care bits (see Table 9). Any additional data bytes sent to the device are ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the RSTE bit in the Status Register is modified, and the WEL bit in the Status Register is reset back to a Logical 0.

The complete one byte of data must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on even byte boundaries (multiples of eight bits); otherwise, the device aborts the operation, the state of the RSTE bit does not change, and the WEL bit in the Status Register is reset back to the Logical 0 state.

Table 9. Write Status Register Byte 2 Format

Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
X	X	X	RSTE	X	X	X	X

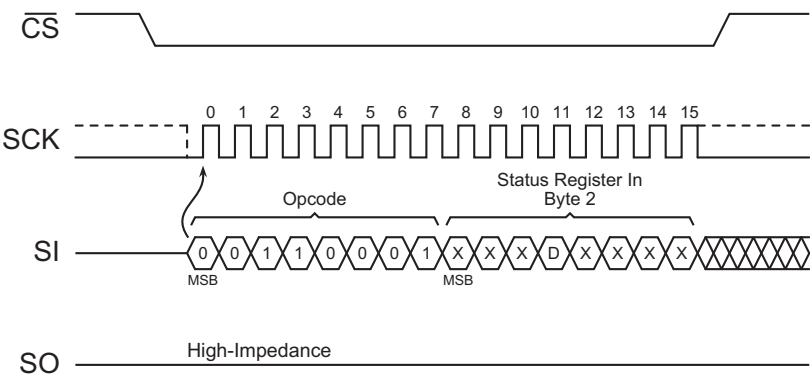


Figure 20. Write Status Register Byte 2

## 12. Other Commands and Functions

### 12.1 Read Manufacturer and Device ID

Identification information can be read from the device to enable systems to electronically query and identify the device while it is in system. The identification method and the command opcode comply with the JEDEC standard for “Manufacturer and Device ID Read Methodology for SPI Compatible Serial Interface Memory Devices”. The type of information that can be read from the device includes the JEDEC defined Manufacturer ID, the vendor specific Device ID, and the vendor specific Extended Device Information.

Since not all Flash devices are capable of operating at very high clock frequencies, applications designed to read the identification information from the devices at a reasonably low clock frequency to ensure all devices used in the application can be identified properly. Once the identification process is complete, the application can increase the clock frequency to accommodate specific Flash devices that are capable of operating at the higher clock frequencies.

To read the identification information, the  $\overline{\text{CS}}$  pin must first be asserted and the opcode of 9Fh must be clocked into the device. After the opcode has been clocked in, the device begins outputting the identification data on the SO pin during the subsequent clock cycles. The first byte to be output is the Manufacturer ID followed by two bytes of Device ID information. The fourth byte output is the Extended Device Information String Length, which is 00h indicating that no Extended Device Information follows. After the Extended Device Information String Length byte is output, the SO pin goes into a high-impedance state; therefore, additional clock cycles have no affect on the SO pin and no data are output. As indicated in the JEDEC standard, reading the Extended Device Information String Length and any subsequent data is optional. Deasserting the  $\overline{\text{CS}}$  pin terminates the Manufacturer and Device ID read operation and put the SO pin into a High-Impedance state. The  $\overline{\text{CS}}$  pin can be deasserted at any time and does not require that a full byte of data be read.

**Table 10. Manufacturer and Device ID Information**

Byte No.	Data Type	Value
1	Manufacturer ID	1Fh
2	Device ID (Part 1)	42h
3	Device ID (Part 2)	00h
4	Extended Device Information String Length	00h

**Table 11. Manufacturer and Device ID Details**

Data Type	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Hex Value	Details
Manufacturer ID	JEDEC Assigned Code								1Fh	JEDEC Code: 0001 1111 (1Fh for Renesas Electronics)
	0	0	0	1	1	1	1	1		
Device ID (Part 1)	Family Code			Density Code					42h	Family Code: 010(AT25F/AT25FSxxx series) Density Code: 00010 (1-Mbit)
	0	1	0	0	0	0	1	0		
Device ID (Part 2)	Sub Code			Product Version Code					00h	Sub Code: 000 (Standard series) Product Version:00000
	0	0	0	0	0	0	0	0		

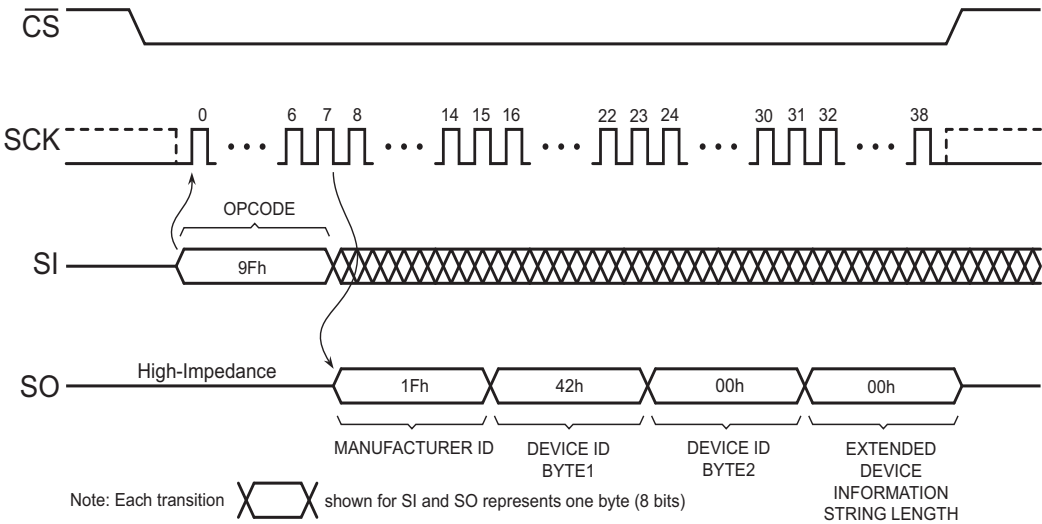


Figure 21. Read Manufacturer and Device ID

## 12.2 Read ID (Legacy Command)

Identification information can be read from the device to enable systems to electronically query and identify the device while it is in system. The preferred method for doing so is the JEDEC standard method described in [“Read Manufacturer and Device ID” on page 31](#); however, the legacy Read ID command is supported on the AT25XE011 to enable backwards compatibility to previous generation devices.

To read the identification information, the  $\overline{CS}$  pin must first be asserted and the opcode of 15h must be clocked into the device. After the opcode has been clocked in, the device begins outputting the identification data on the SO pin during the subsequent clock cycles. The first byte that is output is the Manufacturer ID of 1Fh followed by a single byte of data representing a device code of 65h. After the device code is output, the SO pin goes into a high-impedance state; therefore, additional clock cycles have no affect on the SO pin and no data is output.

Deasserting the  $\overline{CS}$  pin terminates the Read ID operation and puts the SO pin into a high-impedance state. The  $\overline{CS}$  pin can be deasserted at any time and does not require that a full byte of data read.

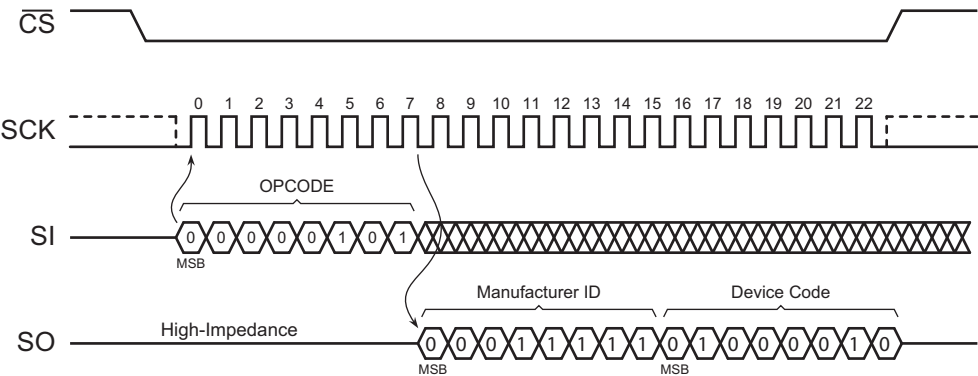


Figure 22. Read ID (Legacy Command)



## 12.3 Deep Power-Down

During normal operation, the device is placed in the standby mode to consume less power as long as the  $\overline{\text{CS}}$  pin remains deasserted and no internal operation is in progress. The Deep Power-Down command offers the ability to place the device into an even lower power consumption state called the Deep Power-Down mode.

When the device is in the Deep Power-Down mode, all commands including the Read Status Register command is ignored with the exception of the Resume from Deep Power-Down command. Since all commands are ignored, the mode can be used as an extra protection mechanism against program and erase operations.

Entering the Deep Power-Down mode is accomplished by simply asserting the  $\overline{\text{CS}}$  pin, clocking in the opcode of B9h, and then deasserting the  $\overline{\text{CS}}$  pin. Any additional data clocked into the device after the opcode is ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the device enters the Deep Power-Down mode within the maximum time of  $t_{\text{EDPD}}$ .

The complete opcode must be clocked in before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise the device aborts the operation and returns to the standby mode once the  $\overline{\text{CS}}$  pin is deasserted. Also, the device defaults to standby mode after a power-cycle.

The Deep Power-Down command is ignored if an internally self-timed operation such as a program or erase cycle is in progress. The Deep Power-Down command must be reissued after the internally self-timed operation has been completed in order for the device to enter the Deep Power-Down mode.

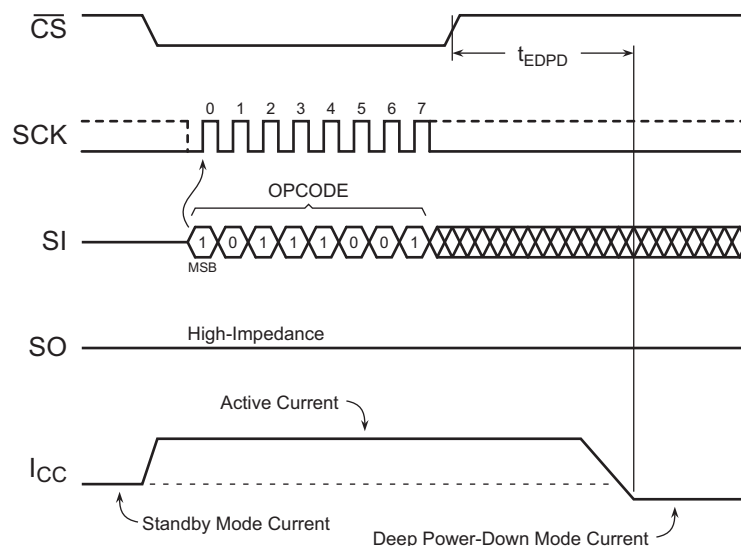


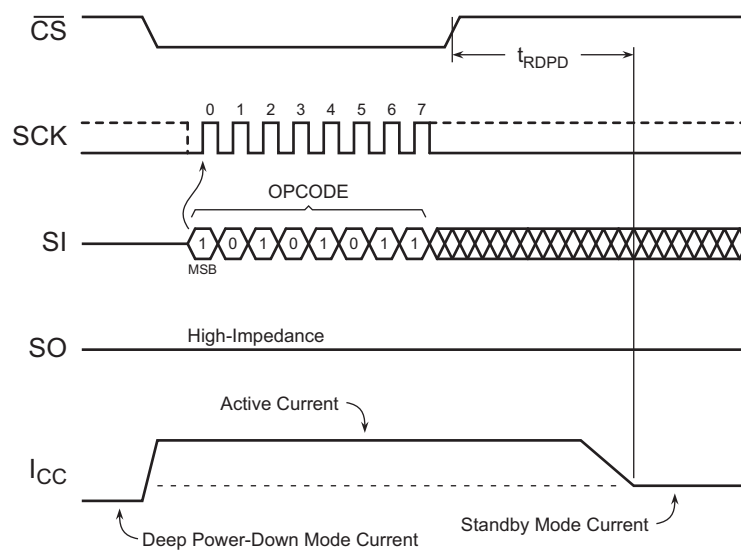
Figure 23. Deep Power-Down

## 12.4 Resume from Deep Power-Down

In order to exit the Deep Power-Down mode and resume normal device operation, the Resume from Deep Power-Down command must be issued. The Resume from Deep Power-Down command is the only command that the device recognizes while in the Deep Power-Down mode.

To resume from the Deep Power-Down mode, the  $\overline{\text{CS}}$  pin must first be asserted and opcode of ABh must be clocked into the device. Any additional data clocked into the device after the opcode is ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the device exits the Deep Power-Down mode within the maximum time of  $t_{\text{RDPD}}$  and return to the standby mode. After the device has returned to the standby mode, normal command operations such as Read Array can be resumed.

If the complete opcode is not clocked in before the  $\overline{\text{CS}}$  pin is deasserted, or if the  $\overline{\text{CS}}$  pin is not deasserted on an even byte boundary (multiples of eight bits), then the device aborts the operation and return to the Deep Power-Down mode.



**Figure 24. Resume From Deep Power-Down**

## 12.5 Ultra-Deep Power-Down

The Ultra-Deep Power-Down mode allows the device to further reduce its energy consumption compared to the existing standby and Deep Power-Down modes by shutting down additional internal circuitry. When the device is in the Ultra-Deep Power-Down mode, all commands including the Status Register Read and Resume from Deep Power-Down commands are ignored. Since all commands are ignored, the mode can be used as an extra protection mechanism against inadvertent or unintentional program and erase operations. Entering the Ultra-Deep Power-Down mode is accomplished by simply asserting the  $\overline{CS}$  pin, clocking in the opcode 79h, and then deasserting the  $\overline{CS}$  pin. Any additional data clocked into the device after the opcode are ignored. When the  $\overline{CS}$  pin is deasserted, the device enters the Ultra-Deep Power-Down mode within the maximum time of  $t_{EUDPD}$ .

The complete opcode must be clocked in before the  $\overline{CS}$  pin is deasserted; otherwise, the device aborts the operation and returns to the standby mode once the  $\overline{CS}$  pin is deasserted. Also, the device defaults to the standby mode after a power cycle. The Ultra-Deep Power-Down command is ignored if an internally self-timed operation such as a program or erase cycle is in progress.

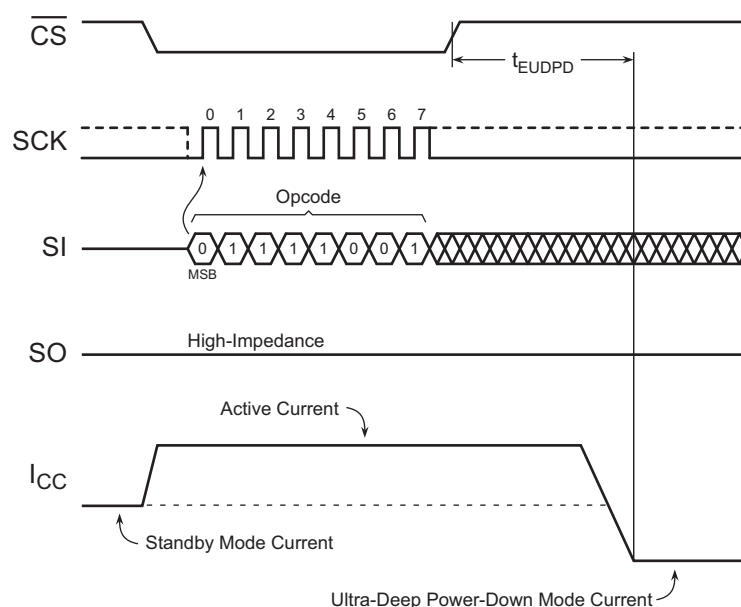


Figure 25. Ultra-Deep Power-Down

## 12.6 Exit Ultra-Deep Power-Down

To exit from the Ultra-Deep Power-Down mode, any one of three operations can be performed:

### Chip Select Toggle

The  $\overline{\text{CS}}$  pin must simply be pulsed by asserting the  $\overline{\text{CS}}$  pin, waiting the minimum necessary  $t_{\text{CSLU}}$  time, and then deasserting the  $\overline{\text{CS}}$  pin again. To facilitate simple software development, a dummy byte opcode can also be entered while the  $\overline{\text{CS}}$  pin is being pulsed; the dummy byte opcode is simply ignored by the device in this case. After the  $\overline{\text{CS}}$  pin has been deasserted, the device exits from the Ultra-Deep Power-Down mode and return to the standby mode within a maximum time of  $t_{\text{XUDPD}}$ . If the  $\overline{\text{CS}}$  pin is reasserted before the  $t_{\text{XUDPD}}$  time has elapsed in an attempt to start a new operation, then that operation is ignored and nothing is performed.

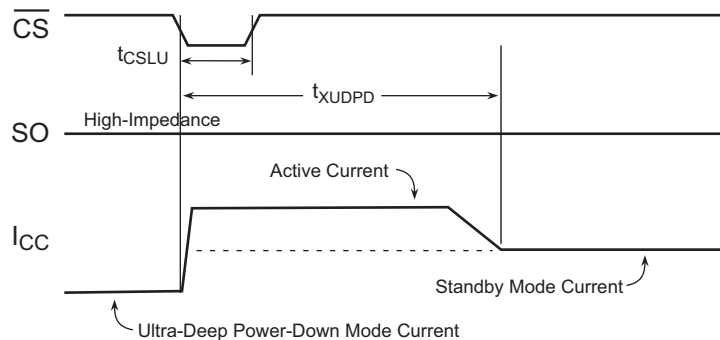


Figure 26. Exit Ultra-Deep Power-Down (Chip Select Toggle)

### Chip Select Low

By asserting the  $\overline{\text{CS}}$  pin, waiting the minimum necessary  $t_{\text{XUDPD}}$  time, and then clocking in the first bit of the next Opcode command cycle. If the first bit of the next command is clocked in before the  $t_{\text{XUDPD}}$  time has elapsed, the device exits Ultra-Deep Power-Down, however the intended operation is ignored.

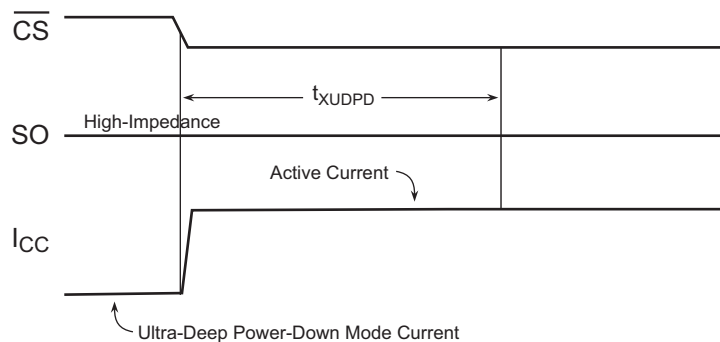


Figure 27. Exit Ultra-Deep Power-Down (Chip Select Low)

### Power Cycling

The device can also exit the Ultra-Deep Power Mode by power cycling the device. The system must wait for the device to return to the standby mode before normal command operations can be resumed. Upon recovery from Ultra-Deep Power-Down all internal registers are at their Power-On default state.

## 12.7 Hold

The  $\overline{\text{HOLD}}$  pin is used to pause the serial communication with the device without having to stop or reset the clock sequence. The Hold mode, however, does not have an effect on any internally self-timed operations such as a program or erase cycle. Therefore, if an erase cycle is in progress, asserting the  $\overline{\text{HOLD}}$  pin does not pause the operation, and the erase cycle continues until it is finished.

The Hold mode can only be entered while the  $\overline{\text{CS}}$  pin is asserted. The Hold mode is activated simply by asserting the  $\overline{\text{HOLD}}$  pin during the SCK low pulse. If the  $\overline{\text{HOLD}}$  pin is asserted during the SCK high pulse, then the Hold mode won't be started until the beginning of the next SCK low pulse. The device remains in the Hold mode as long as the  $\overline{\text{HOLD}}$  pin and  $\overline{\text{CS}}$  pin are asserted.

While in the Hold mode, the SO pin is in a high-impedance state. Also, both the SI pin and the SCK pin are ignored. The  $\overline{\text{WP}}$  pin, however, can still be asserted or deasserted while in the Hold mode.

To end the Hold mode and resume serial communication, the  $\overline{\text{HOLD}}$  pin must be deasserted during the SCK low pulse. If the  $\overline{\text{HOLD}}$  pin is deasserted during the SCK high pulse, then the Hold mode won't end until the beginning of the next SCK low pulse.

If the  $\overline{\text{CS}}$  pin is deasserted while the  $\overline{\text{HOLD}}$  pin is still asserted, then any operation that has been started is aborted, and the device resets the WEL bit in the Status Register back to the logical 0 state.

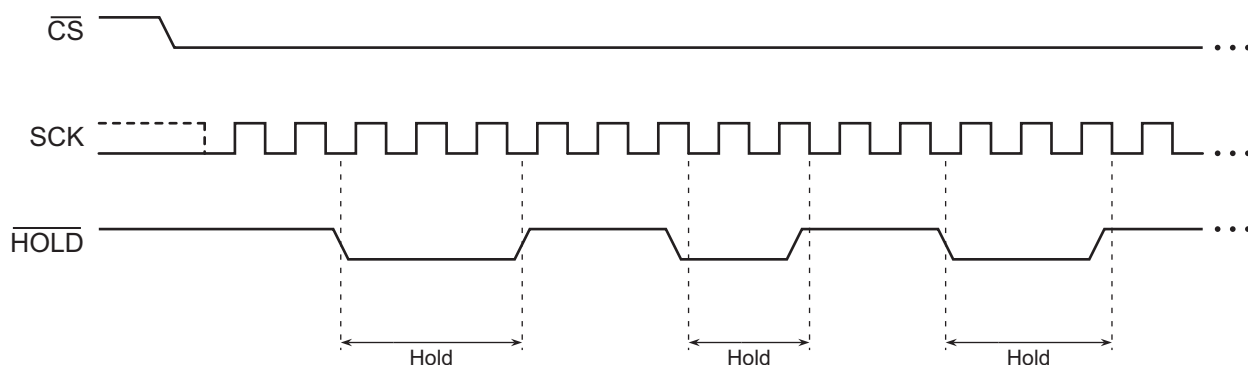


Figure 28. Hold Mode

## 12.8 Reset

In some applications, it can be necessary to prematurely terminate a program or erase operation rather than wait the hundreds of microseconds or milliseconds necessary for the program or erase operation to complete normally. The Reset command allows a program or erase operation in progress to be ended abruptly and returns the device to an idle state. Since the need to reset the device is immediate, the Write Enable command does not need to be issued prior to the Reset command. Therefore, the Reset command operates independently of the state of the WEL bit in the Status Register.

The Reset command can be executed only if the command has been enabled by setting the Reset Enabled (RSTE) bit in the Status Register to a Logical 1 using write status register byte 2 command 31h. This command must be entered before a program command is entered. If the Reset command has not been enabled (the RSTE bit is in the Logical 0 state), then any attempts at executing the Reset command is ignored.

To perform a Reset, the  $\overline{\text{CS}}$  pin must first be asserted, and then the opcode F0h must be clocked into the device. No address bytes need to be clocked in, but a confirmation byte of D0h must be clocked into the device immediately after the opcode. Any additional data clocked into the device after the confirmation byte is ignored. When the  $\overline{\text{CS}}$  pin is deasserted, the program operation currently in progress is terminated within a time of  $t_{\text{SWRST}}$ . Since the program or erase operation may not complete before the device is reset, the contents of the page being programmed or erased cannot be guaranteed to be valid.

The Reset command has no effect on the states of the Configuration Register or RSTE bit in the Status Register. The WEL however, is reset back to its default state.

The complete opcode and confirmation byte must be clocked into the device before the  $\overline{\text{CS}}$  pin is deasserted, and the  $\overline{\text{CS}}$  pin must be deasserted on an even byte boundary (multiples of eight bits); otherwise, no Reset operation is performed.

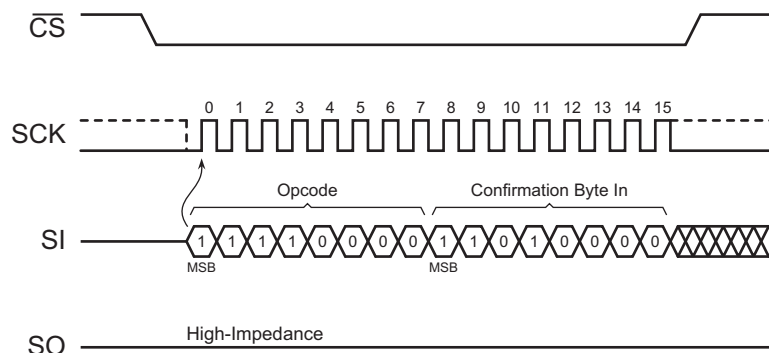


Figure 29. Reset

## 13. Electrical Specifications

Temperature under Bias . . . . . -55 °C to +125 °C	<p><b>*Notice:</b> Stresses beyond those listed under “Absolute Maximum Ratings” can cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions beyond those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods can affect device reliability. Voltage extremes referenced in the “Absolute Maximum Ratings” are intended to accommodate short duration undershoot/overshoot conditions and does not imply or guarantee functional device operation at these levels for any extended period of time.</p>
Storage Temperature . . . . . -65 °C to +150 °C	
Absolute Maximum $V_{CC}$ . . . . . 3.96 V	
All Output Voltages with Respect to Ground . . . . . -0.6 V to 4.2 V (Max $V_{CC}$ of 3.6 V + 0.6 V)	
All Output Voltages with Respect to Ground (excluding $V_{CC}$ pin, including NC pins) . . . . . -0.6 V to 4.2 V (Max $V_{CC}$ of 3.6 V + 0.6 V)	

### 13.1 DC and AC Operating Range

Parameter		AT25XE011
Operating Temperature (Case) <sup>[1]</sup>	Industrial	-40 °C to 85 °C
$V_{CC}$ Power Supply		1.65 V to 3.6 V

1. Temperature Range: -10 °C to +85 °C (1.65V to 3.6V), -40 °C to +85 °C (1.7 V to 3.6 V).

## 13.2 DC Characteristics

Symbol	Parameter	Condition	1.65V to 3.6V			2.3V to 3.6V			Units
			Min	Typ	Max	Min	Typ	Max	
$I_{UDPD}$	Ultra-Deep Power-Down Current	$\overline{CS} = V_{CC}$ . All other inputs at 0 V or $V_{CC}$		0.2	0.6		0.3	0.6	$\mu A$
$I_{DPD}^{[1]}$	Deep Power-Down Current	$\overline{CS} = V_{CC}$ . All other inputs at 0 V or $V_{CC}$		4.5	15		5	15	$\mu A$
$I_{SB}$	Standby Current	$\overline{CS} = V_{CC}$ . All other inputs at 0 V or $V_{CC}$		25	40		25	40	$\mu A$
$I_{CC1}^{[2][3]}$	Active Current, Low Power Read (03h, 0Bh) Operation	$f = 1 \text{ MHz}$ ; $I_{OUT} = 0 \text{ mA}$		3.5	7		5	7	mA
		$f = 20 \text{ MHz}$ ; $I_{OUT} = 0 \text{ mA}$		3.5	7		6	7	mA
$I_{CC2}^{[2][3]}$	Active Current, Read Operation	$f = 50 \text{ MHz}$ ; $I_{OUT} = 0 \text{ mA}$		4	7		5.5	7	mA
		$f = 85 \text{ MHz}$ ; $I_{OUT} = 0 \text{ mA}$		4	7		5.5	7	mA
$I_{CC3}^{[2][3]}$	Active Current, Program Operation	$\overline{CS} = V_{CC}$		10	13		10	13	mA
$I_{CC4}^{[2][3]}$	Active Current, Erase Operation	$\overline{CS} = V_{CC}$		9	13		10	13	mA
$I_{LI}$	Input Load Current	All inputs at CMOS levels			1			1	$\mu A$
$I_{LO}$	Output Leakage Current	All inputs at CMOS levels			1			1	$\mu A$
$V_{IL}$	Input Low Voltage				$V_{CC} \times 0.2$			$V_{CC} \times 0.3$	V
$V_{IH}$	Input High Voltage		$V_{CC} \times 0.8$			$V_{CC} \times 0.7$			V
$V_{OL}$	Output Low Voltage	$I_{OL} = 100 \mu A$			0.2			0.4	V
$V_{OH}$	Output High Voltage	$I_{OH} = -100 \mu A$	$V_{CC} - 0.2V$			$V_{CC} - 0.2V$			V

1. Max. specification is 20  $\mu A$  @ 85 °C.

2. Typical values measured at 1.8 V @ 25 °C for the 1.65 V to 3.6 V range.

3. Typical values measured at 3.0 V @ 25 °C for the 2.3 V to 3.6 V range.

## 13.3 AC Characteristics - Maximum Clock Frequencies

Symbol	Parameter	1.65V to 3.6V			2.3V to 3.6V			Units
		Min	Typ	Max	Min	Typ	Max	
$f_{CLK}$	Maximum Clock Frequency for All Operations (including 0Bh opcode)			104			104	MHz
$f_{RDLF}$	Maximum Clock Frequency for 03h Opcode (Read Array – Low Frequency)			25			33	MHz
$f_{RDDO}$	Maximum Clock Frequency for 3B Opcode			50			50	MHz



## 13.4 AC Characteristics - All Other Parameters

Symbol	Parameter	1.65V to 3.6V			2.3V to 3.6V			Units
		Min	Typ	Max	Min	Typ	Max	
$t_{CLKH}$	Clock High Time	4			4			ns
$t_{CLKL}$	Clock Low Time	4			4			ns
$t_{CLKR}^{[1]}$	Clock Rise Time, Peak-to-Peak (Slew Rate)	0.1			0.1			V/ns
$t_{CLKF}^{[1]}$	Clock Fall Time, Peak-to-Peak (Slew Rate)	0.1			0.1			V/ns
$t_{CSH}$	Chip Select High Time	35			25			ns
$t_{CSLS}$	Chip Select Low Setup Time (relative to Clock)	6			6			ns
$t_{CSLH}$	Chip Select Low Hold Time (relative to Clock)	6			6			ns
$t_{CSHS}$	Chip Select High Setup Time (relative to Clock)	6			6			ns
$t_{CSHH}$	Chip Select High Hold Time (relative to Clock)	6			6			ns
$t_{DS}$	Data In Setup Time	2			2			ns
$t_{DH}$	Data In Hold Time	1			1			ns
$t_{DIS}^{[1]}$	Output Disable Time			8			6	ns
$t_V$	Output Valid Time			8			6	ns
$t_{OH}$	Output Hold Time	0			0			ns
$t_{HLS}$	$\overline{HOLD}$ Low Setup Time (relative to Clock)	6			5			ns
$t_{HLH}$	$\overline{HOLD}$ Low Hold Time (relative to Clock)	6			5			ns
$t_{HHS}$	$\overline{HOLD}$ High Setup Time (relative to Clock)	6			5			ns
$t_{HHH}$	$\overline{HOLD}$ High Hold Time (relative to Clock)	6			5			ns
$t_{HLQZ}^{[1]}$	$\overline{HOLD}$ Low to Output High-Z			7			6	ns
$t_{HHQX}^{[1]}$	$\overline{HOLD}$ High to Output Low-Z			7			6	ns
$t_{WPS}^{[1][2]}$	Write Protect Setup Time	20			20			ns
$t_{WPH}^{[1][2]}$	Write Protect Hold Time	100			100			ns
$t_{EDPD}^{[1]}$	Chip Select High to Deep Power-Down			2			2	$\mu$ s
$t_{EUDPD}$	Chip Select High to Ultra-Deep Power-Down			3			3	$\mu$ s
$t_{SWRST}$	Software Reset Time			60			60	$\mu$ s
$t_{CSLU}$	Minimum Chip Select Low to Exit Ultra-Deep Power-Down	20			20			ns
$t_{XUDPD}$	Exit Ultra-Deep Power-Down Time	70			70			$\mu$ s
$t_{RDPD}^{[1]}$	Chip Select High to Standby Mode			8			8	$\mu$ s

1. Not 100% tested (value guaranteed by design and characterization).
2. Only applicable as a constraint for the Write Status Register command when BPL = 1.

### 13.5 Program and Erase Characteristics

Symbol	Parameter		1.65V-3.6V			2.3V-3.6V			
			Min	Typ	Max	Min	Typ	Max	Units
t <sub>PP</sub> <sup>[1]</sup>	Page Program Time	256 Bytes		2	3		2	3	ms
t <sub>BP</sub>	Byte Program Time			12			8		μs
t <sub>PE</sub>	Page Erase Time	256 Bytes		7	25		7	25	ms
t <sub>BLKE</sub> <sup>[1]</sup>	Block Erase Time	4 kBytes		50	75		50	75	ms
		32 kBytes		400	500		380	450	
t <sub>CHPE</sub> <sup>[1][2]</sup>	Chip Erase Time			1600	2200		1600	2000	ms
t <sub>OTPP</sub>	OTP Security Register Program Time			400	950		400	950	μs
t <sub>WRSR</sub> <sup>[2]</sup>	Write Status Register Time			20	40		20	40	ms

1. Maximum values indicate worst-case performance after 100,000 erase/program cycles.
2. Not 100% tested (value guaranteed by design and characterization).
3. Program and Erase operations characterized at -10 °C to +85 °C. Program and Erase operations at -40 °C to -10 °C require a minimum of 1.7 V.

## 14. Power-On/Reset State

When power is first applied to the device, or when recovering from a reset condition, the output pin (SO) is in a high impedance state, and a high-to-low transition on the CSB pin is required to start a valid instruction. The SPI mode (Mode 3 or Mode 0) is automatically selected on every falling edge of CSB by sampling the inactive clock state.

### 14.1 Power-Up/Power-Down Voltage and Timing Requirements

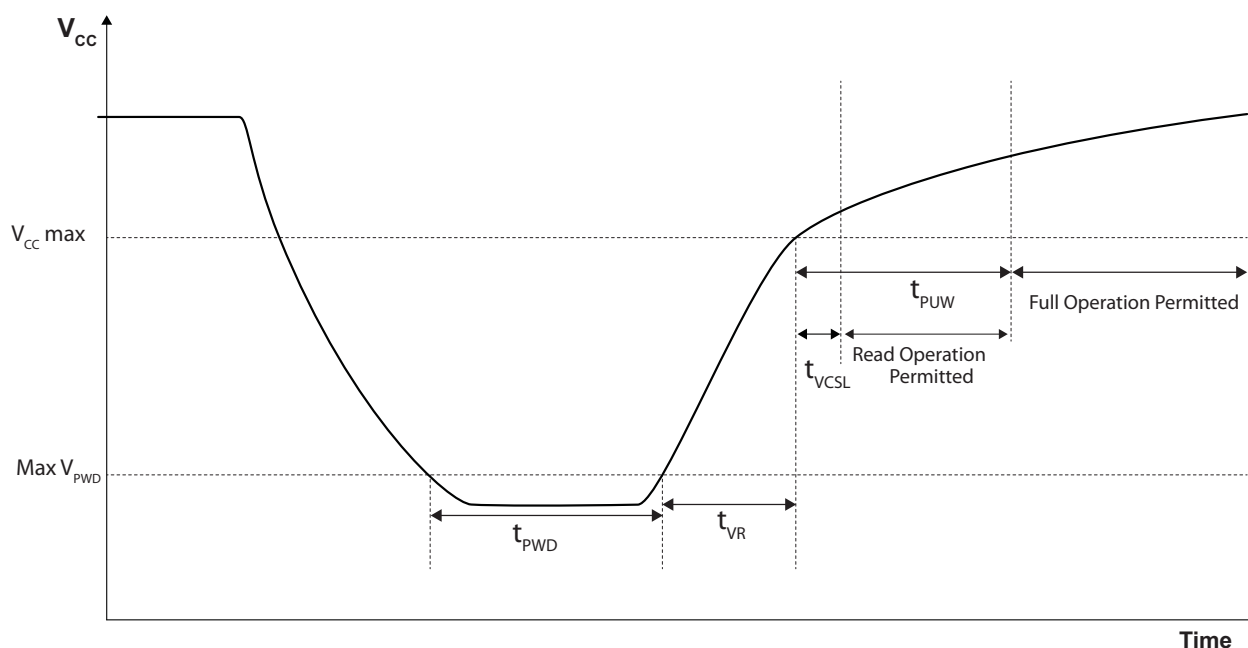
During power-up, the device must not be read for at least the minimum  $t_{VCSL}$  time after the supply voltage reaches the minimum  $V_{POR}$  level ( $V_{POR\ min}$ ). While the device is being powered-up, the internal Power-On Reset (POR) circuitry keeps the device in a reset mode until the supply voltage rises above the minimum  $V_{CC}$ . During this time, all operations are disabled and the device does not respond to any commands.

If the first operation to the device after power-up is a program or erase operation, the operation cannot be started until the supply voltage reaches the minimum  $V_{CC}$  level and an internal device delay has elapsed. This delay has a maximum time of  $t_{PUW}$ . After  $t_{PUW}$  time, the device is in standby mode if CSB is at logic high, or active mode if CSB is at logic low. For the case of Power-down then Power-up operation, or if a power interruption occurs (such that  $V_{CC}$  drops below  $V_{POR\ max}$ ), the  $V_{CC}$  of the Flash device must be maintained below  $V_{PWD}$  for at least the minimum specified  $T_{PWD}$  time. This ensures the Flash device resets properly after a power interruption.

**Table 12. Voltage and Timing Requirements for Power-Up/Power-Down**

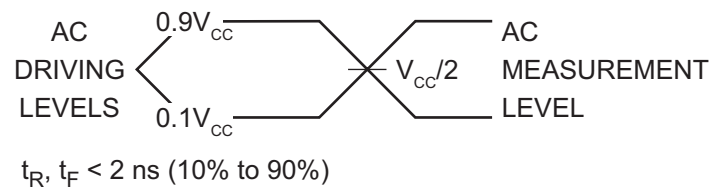
Symbol	Parameter	Min	Max	Units
$V_{PWD}^{[1]}$	$V_{CC}$ for device initialization		1.0	V
$t_{PWD}^{[1]}$	Minimum duration for device initialization	300		$\mu s$
$t_{VCSL}$	Minimum $V_{CC}$ to chip select low time for Read command	70		$\mu s$
$t_{VR}^{[1]}$	$V_{CC}$ rise time	1	500000	$\mu s/V$
$V_{POR}$	Power on reset voltage	1.45	1.6	V
$t_{PUW}$	Power up delay time before Program or Erase is allowed		3	ms

1. Not 100% tested (value guaranteed by design and characterization).

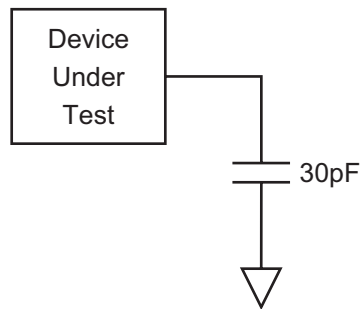


**Figure 30. Power-Up Timing**

## 14.2 Input Test Waveforms and Measurement Levels



## 14.3 Output Test Load



15. AC Waveforms

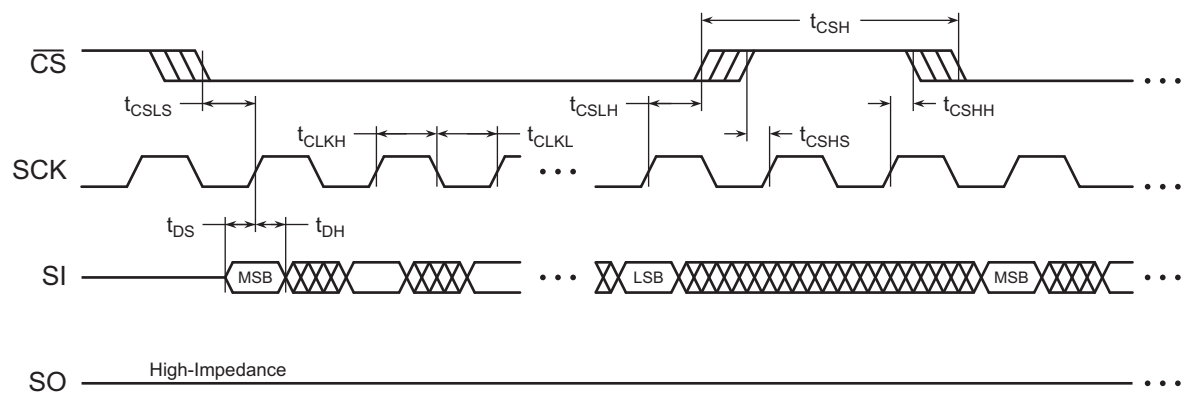


Figure 31. Serial Input Timing

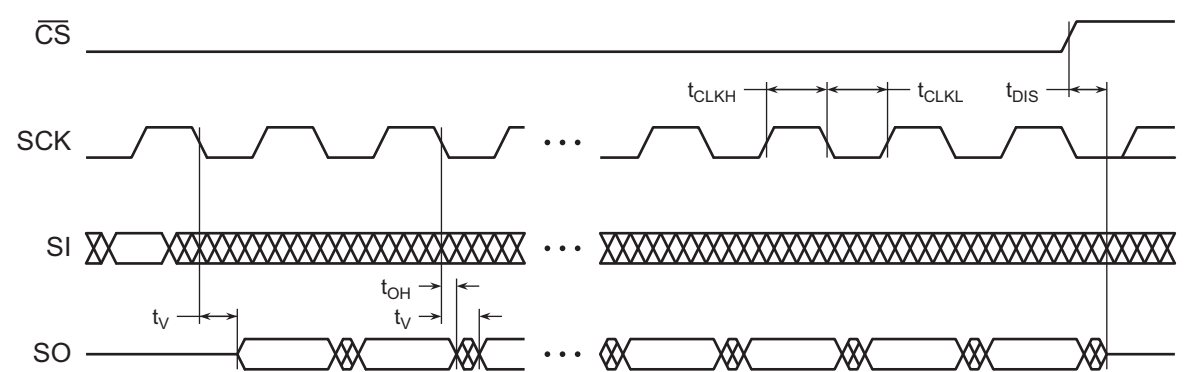


Figure 32. Serial Output Timing

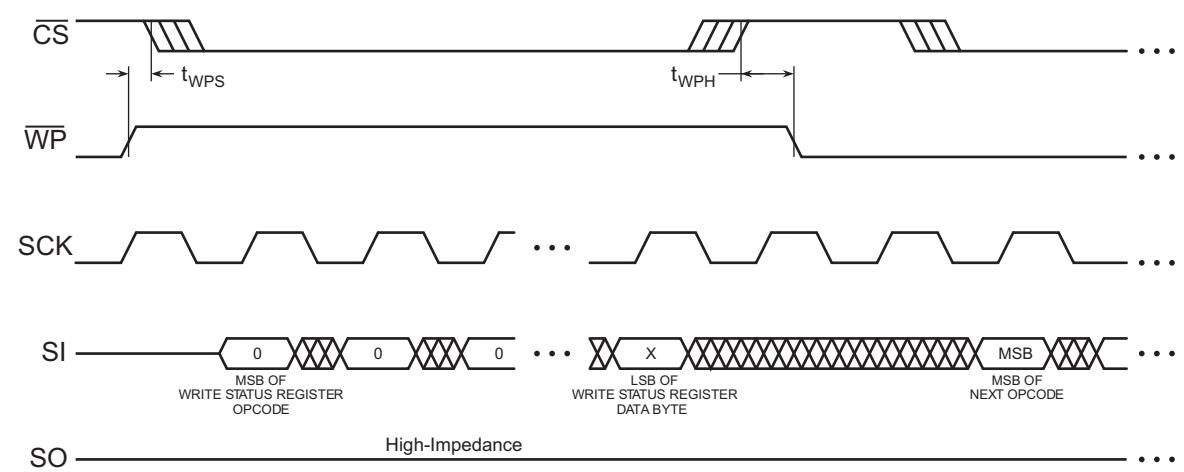


Figure 33.  $\overline{WP}$  Timing for Write Status Register Command When BPL=1

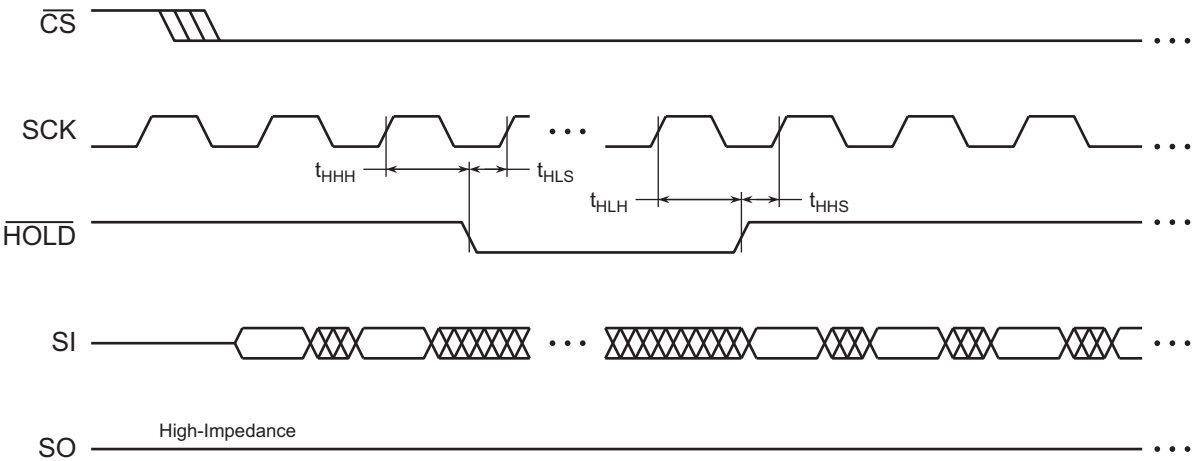


Figure 34.  $\overline{\text{HOLD}}$  Timing - Serial Input

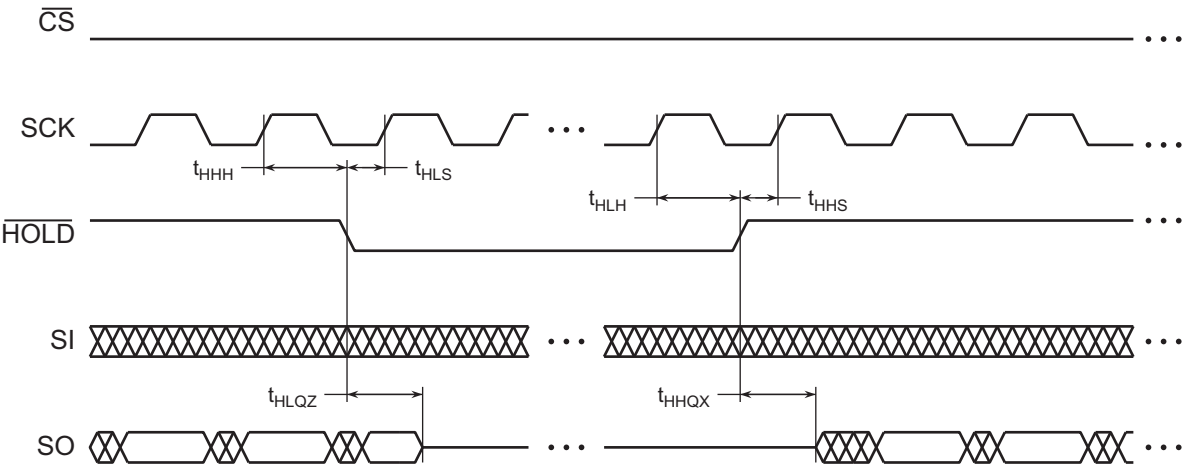
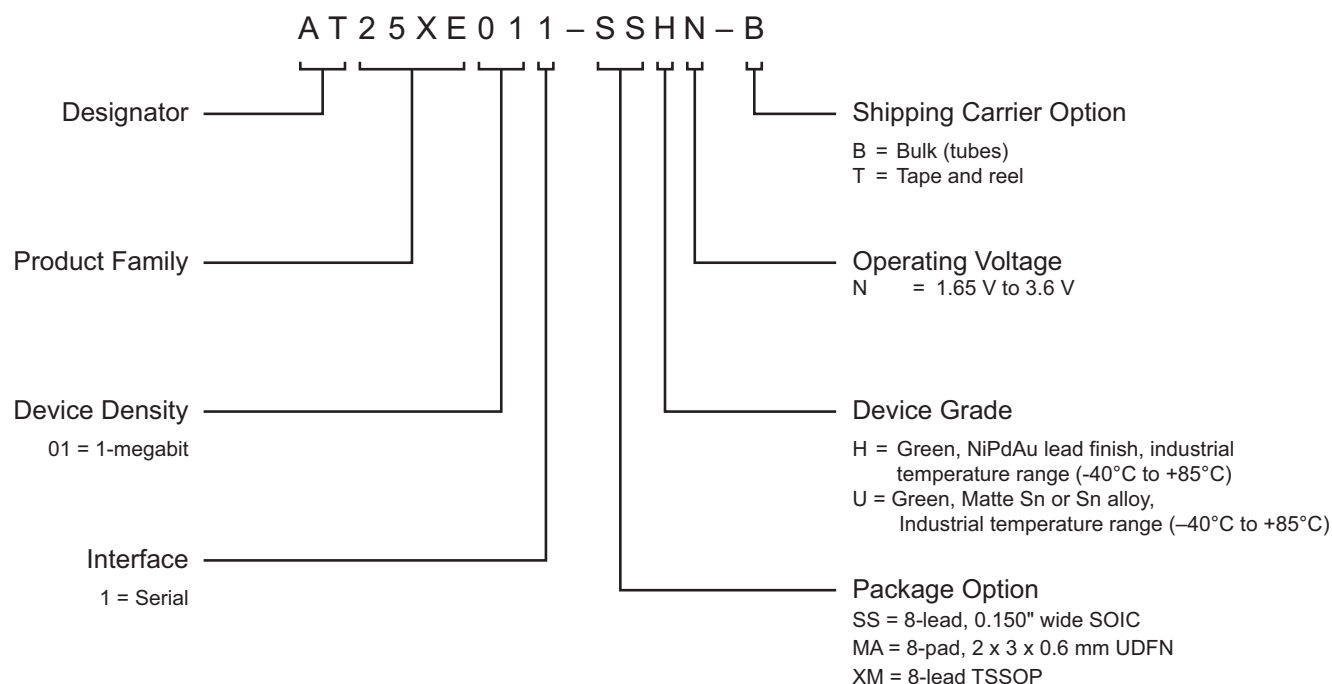


Figure 35.  $\overline{\text{HOLD}}$  Timing - Serial Output

## 16. Ordering Information

### 16.1 Ordering Code Detail



Ordering Code <sup>[1]</sup>	Package	Lead Finish	Operating Voltage	Max. Freq. (MHz)	Operation Range
AT25XE011-SSH-N-B AT25XE011-SSH-N-T	8S1	NiPdAu	1.65 V to 3.6 V <sup>[2]</sup>	104	Industrial (–40°C to +85°C) <sup>[2]</sup>
AT25XE011-MAHN-T	8MA3				
AT25XE011-XMHN-B AT25XE011-XMHN-T	8X				

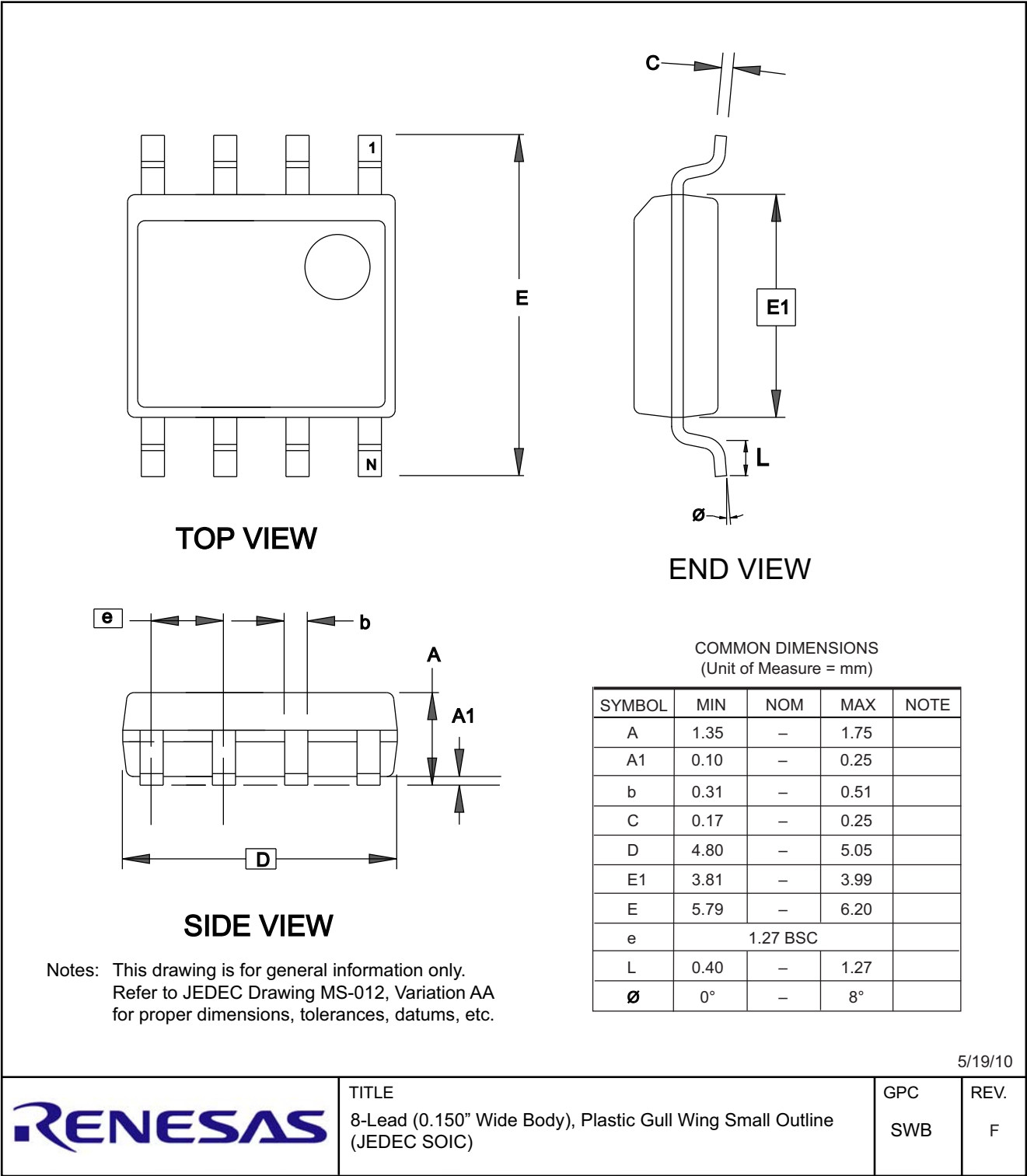
1. The shipping carrier option code is not marked on the devices.

2. Temperature Range: –10°C to +85°C (1.65V to 3.6V), –40°C to +85°C (1.7V to 3.6V).

Package Type	
8S1	8-lead, 0.150" Wide, Plastic Gull Wing Small Outline Package (JEDEC SOIC)
8MA3	8-pad, 2 x 3 x 0.6 mm, Thermally Enhanced Plastic Ultra Thin Dual Flat No Lead Package (UDFN)
8X	8-lead, Thin Small Outline Package

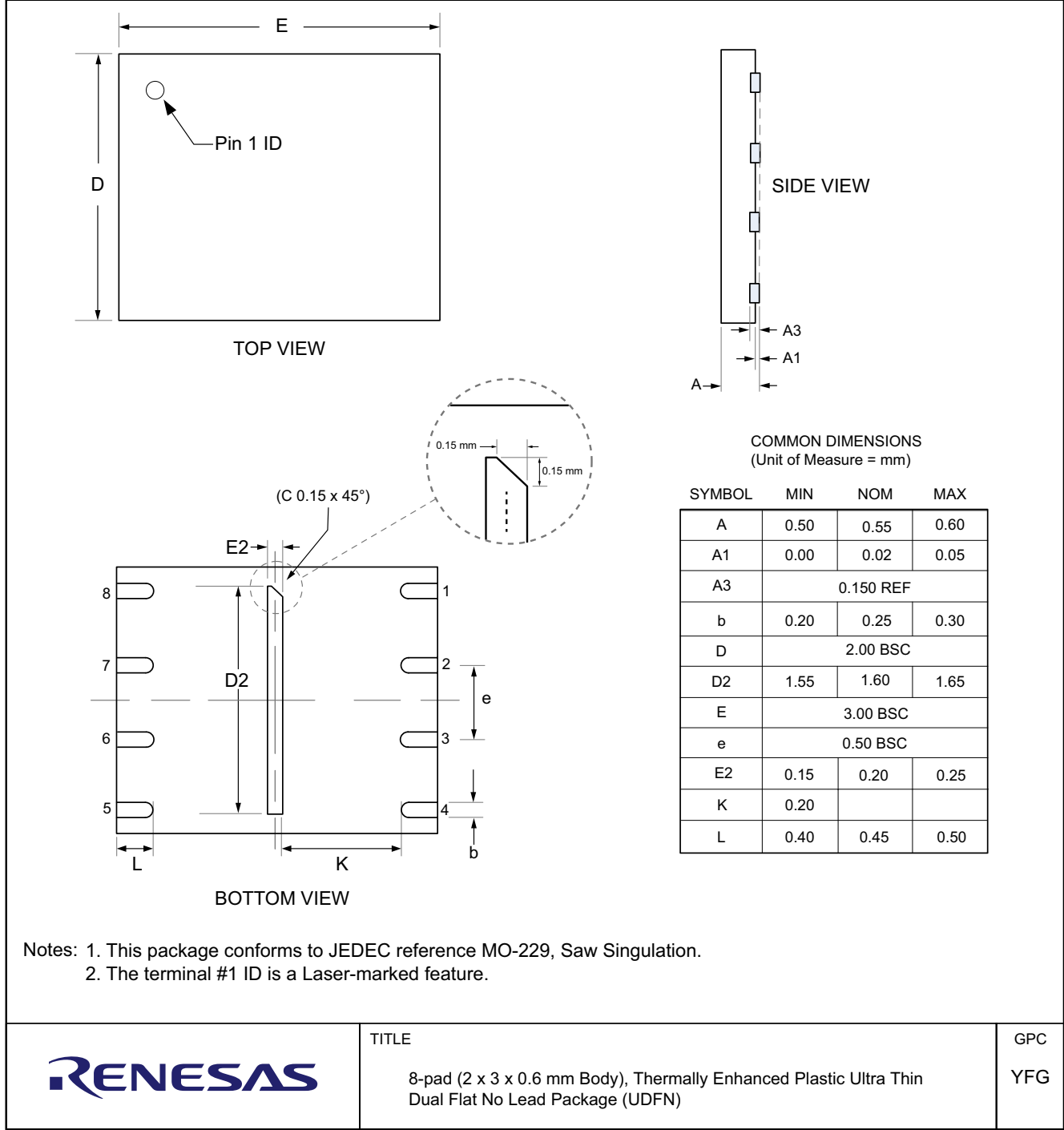
17. Packaging Information

17.1 8S1 – JEDEC SOIC

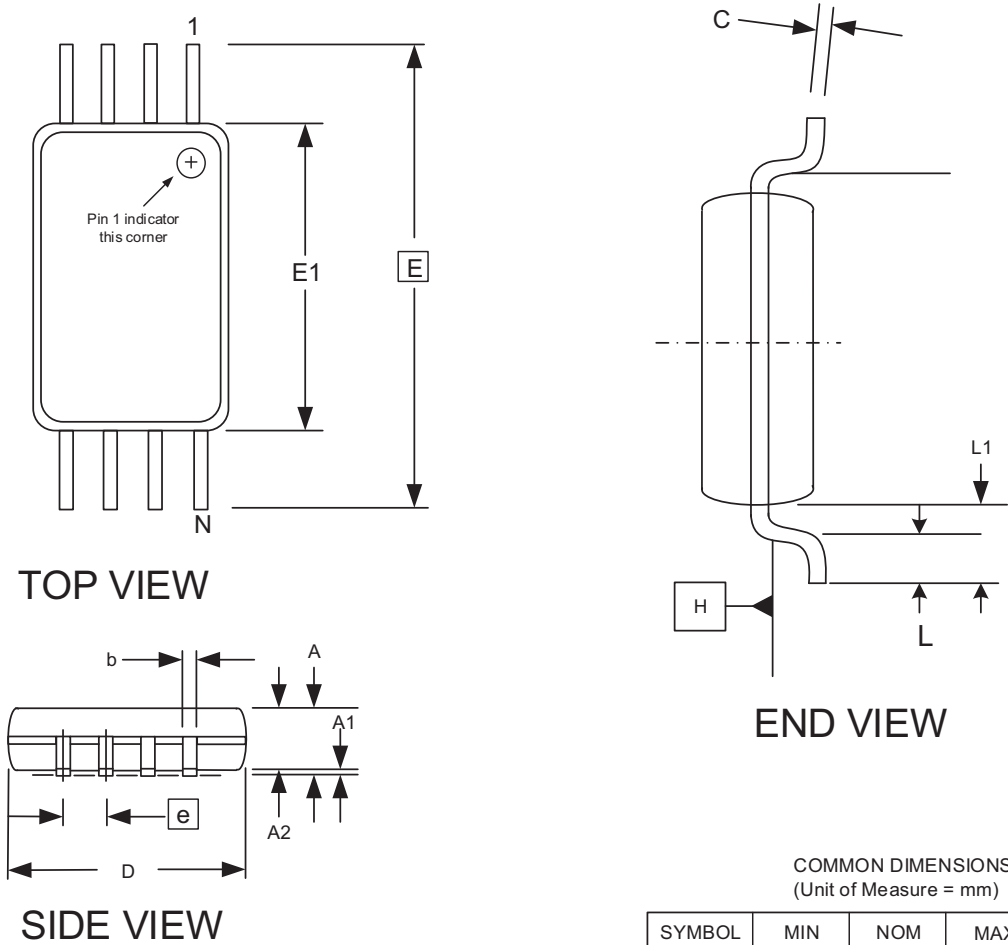




17.2 8MA3 UDFN



17.3 8X-TSSOP



- Notes: 1. This drawing is for general information only. Refer to JEDEC Drawing MO-153. Variation AA, for proper dimensions, tolerances, datums, etc.
2. Dimension D does not include mold Flash, protrusions or gate burrs. Mold Flash, protrusions and gate burrs shall not exceed 0.15mm (0.006 in) per side.
3. Dimension E1 does not include inter-lead Flash or protrusions, Inter-lead Flash and protrusions shall not exceed 0.2 5mm (0.01 in) per side.
4. Dimension b does not include Dambar protrusion. Allowable Dambar protrusion shall be 0.08 mm total in excess of the b dimension at maximum material condition. Dambar cannot be located on the lower radius of the foot. Minimum space between protrusion and adjacent lead is 0.07 mm.
5. Dimension D and E1 to be determined at Datum Plane H.

COMMON DIMENSIONS  
(Unit of Measure = mm)

SYMBOL	MIN	NOM	MAX	NOTE
A	--	—	1.20	
A1	0.05	—	0.15	
A2	0.80	1.00	1.05	
D	2.90	3.00	3.10	2, 5
E	6.40 BSC			
E1	4.30	4.40	4.50	3, 5
b	0.19	—	0.30	4
e	0.65 BSC			
L	0.45	0.60	0.75	
L1	1.00 REF			
C	0.09	—	0.20	

	TITLE 8X, 8-lead 4.4mm Body. Plastic Thin Shrink Small Outline Package (TSSOP)	GPC	DRAWING NO.	REV.
		TNR	8X	E

## 18. Revision History

Revision Level – Release Date	History
A – October 2014	Initial release.
B – October 2014	Document status changed to Preliminary.
C – February 2015	Added IDPD specification footnote.
D – November 2015	Updated IDPD and ISB specification conditions. Updated document status from Preliminary to Complete.
E – September 2016	Added WLCSP package. Added U device grade.
F – February 2017	Added patent information. Updated description in Section 8.2 (Page Erase). Updated Power-On Timing description (Section 14.)
G - June 2017	Added clarification of Absolute Maximum Ratings. Added footnote and updated WLCSP references.
H - December 2021	Updated package drawing in section 17.2.
I – March 2022	Changed company logo to Renesas. Removed references to 8-ball WLCSP in document.
J - August 2022	Applied new corporate template to document.